

COMP 250

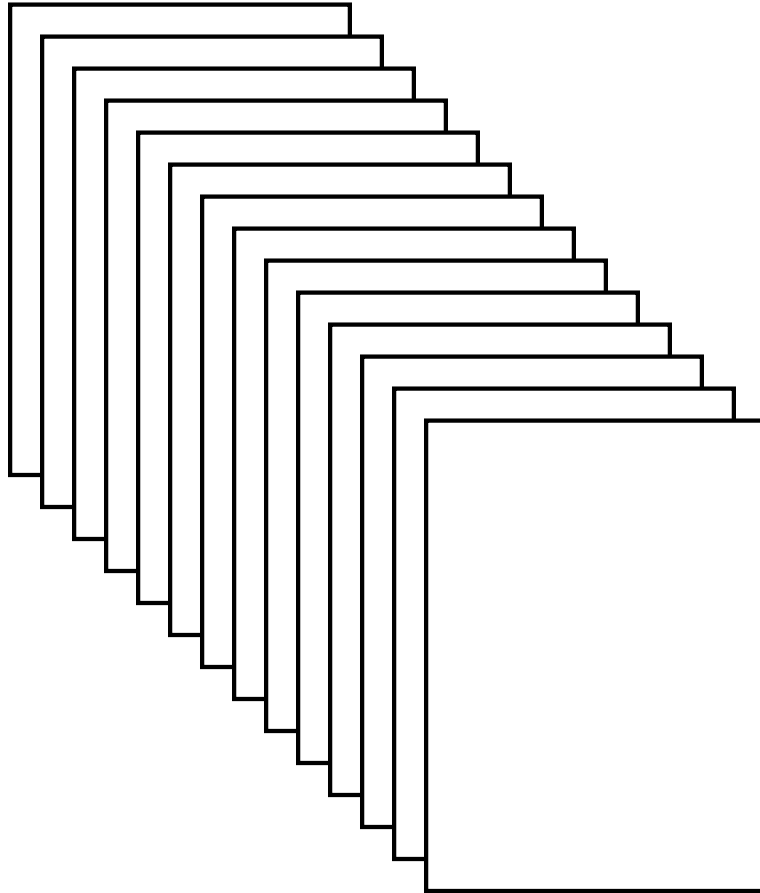
Lecture 12

Algorithms for Sorting a List:

bubble sort  
selection sort  
insertion sort

Feb. 2, 2022

# Example 1: sorting exams by last name



## Example 2: Email packets

When you send a large file by email, it gets broken down into small pieces called “packets” and each packet takes an independent network path to the destination.

Then the packets need to be put together again in their correct order.

<https://computer.howstuffworks.com/question525.htm>

Some sorting algorithms are faster than others.  
See visualization:

<https://www.youtube.com/watch?v=ZZuD6iUe3Pc>

Barak Obama knows about sorting...



[https://www.youtube.com/watch?v=k4RRi\\_ntQc8](https://www.youtube.com/watch?v=k4RRi_ntQc8)

# Sorting Algorithms

- Bubble sort
  - Selection sort
  - Insertion sort
- } today  $O(N^2)$

- Mergesort
  - Heapsort
  - Quicksort
- } later  $O(N \log N)$

# Sorting Algorithms

Today we are concerned with algorithms, not data structures.

Today's algorithms can be implemented easily using an array list or a (doubly) linked list.

# Notation for today...

BEFORE

0	3
1	17
2	-5
3	-2
4	23
5	4

sort into  
increasing  
order

AFTER

-5
-2
3
4
17
23

# Bubble Sort

Given a list of size  $N$ , arrange the elements in *increasing* order.

Pass through the list  $N$  times.

For each pass,

if two neighboring elements are in the wrong order,  
then swap them.



The name invokes the (vague)  
metaphor of bubbles rising in a liquid.



# Bubble Sort Algorithm

```
for  $i = 0$  to  $N - 1$  {           // i-th pass
    for  $k = 0$  to  $N - 2$  {
        if ( list[k] > list[k+1] ) { // wrong order
            list.swap(k, k+1)
        }
    }
}
```

# Example: first pass

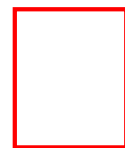
0	3
1	17
2	-5
3	23
4	-2
5	4

```
if list[ 0 ] > list[ 1 ]           // wrong order
    swap( list[ 0 ], list[1 ] )
```

# Example: first pass

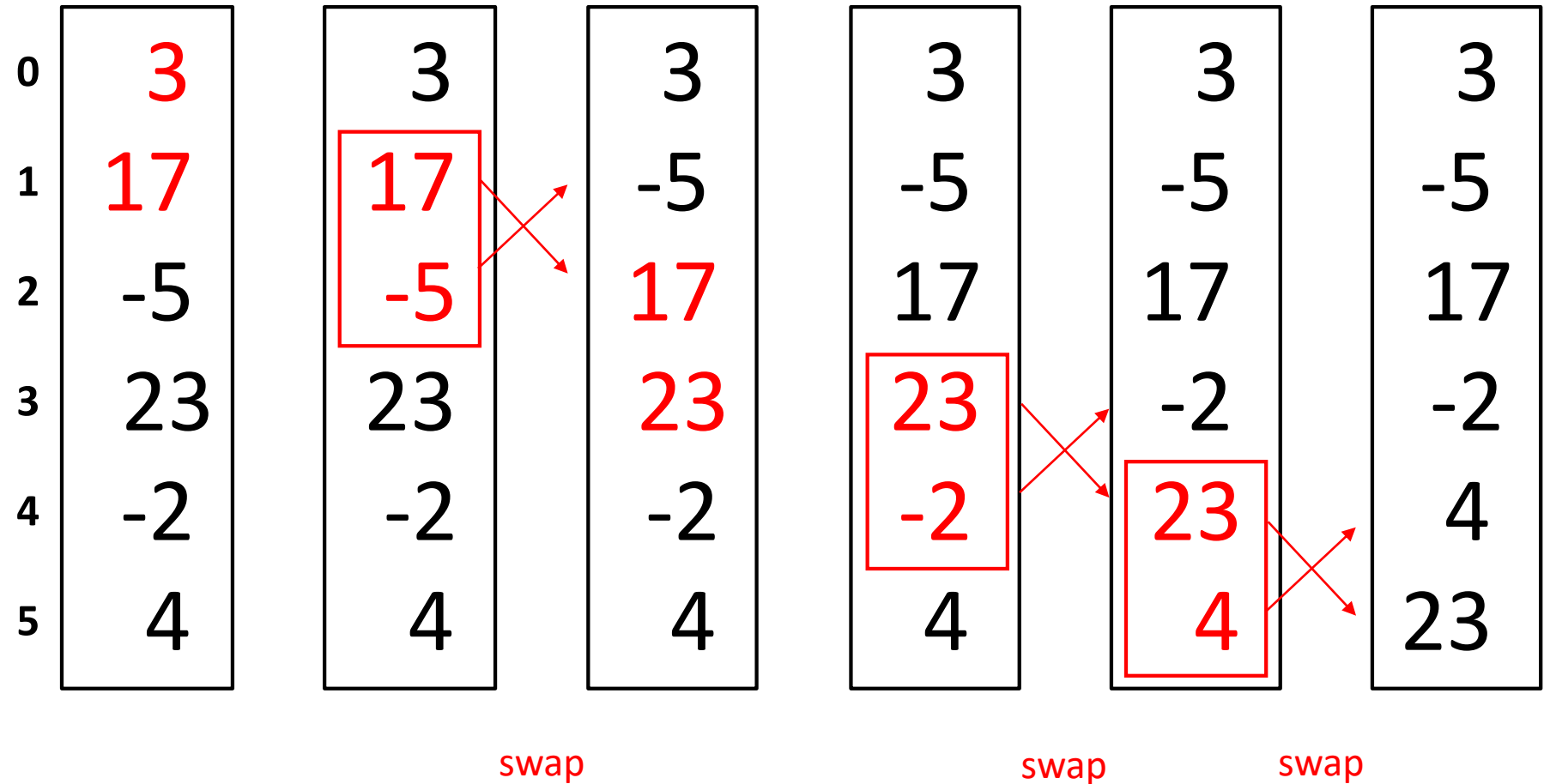
0	3	3
1	17	17
2	-5	-5
3	23	23
4	-2	-2
5	4	4

```
if list[ 1 ] > list[ 2 ]  
    list.swap( 1, 2 )
```

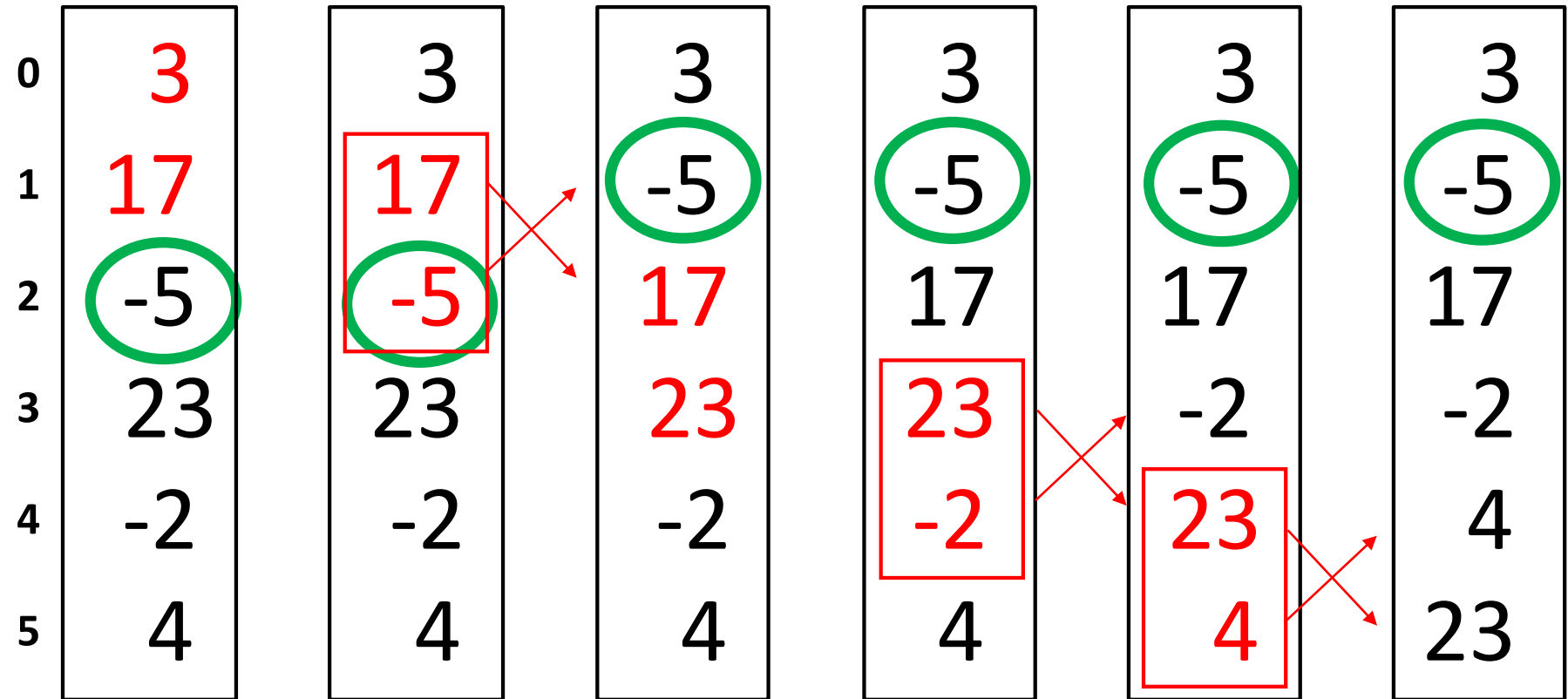


Indicates  
elements need  
to be swapped

# Example: first pass

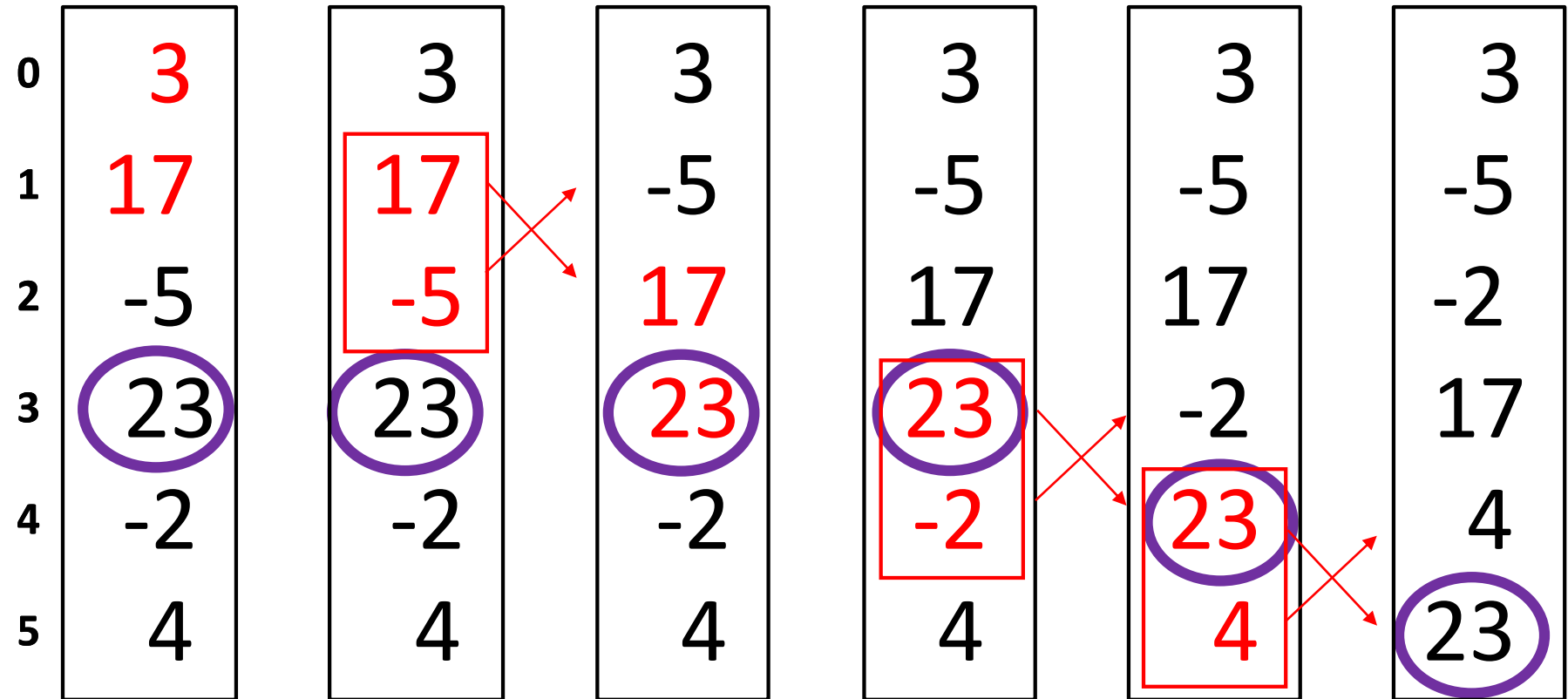


# Smallest element moves up\*



\*assuming it wasn't already at the front of the list

# Largest element moves down \*



\*assuming it wasn't already at the end of the list

What can we say at end of the first pass?

Q: Where is the largest element ?

A: It must be at the end of the list (position  $N-1$ ).

Q: Where is the smallest element ?

A: Could be anywhere except position  $N-1$ .

# Bubble Sort Algorithm

```
for  $i = 0$  to  $N - 1$  {  
    for  $k = 0$  to  $N - 2 - i$  {  
        if ( list[ $k$ ] > list[ $k+1$ ] ) {  
            list.swap(  $k$ ,  $k+1$  )  
        }  
    }  
}
```

Before pass  $i$ , the largest  $i$  elements must already be in their correct position at the end of the list.

Thus, the inner loop can get shorter each time.



# Bubble Sort Algorithm

```
for  $i = 0$  to  $N - 2$  {  
    for  $k = 0$  to  $N - 2 - i$  {  
        if ( list[ $k$ ] > list[ $k+1$ ] ) {  
            list.swap(  $k$ ,  $k+1$  )  
        }  
    }  
}
```

The outer loop only needs to run  $N - 1$  times.

(If the largest  $N - 1$  elements are in their correct position, then the smallest element must also be in its correct position.)

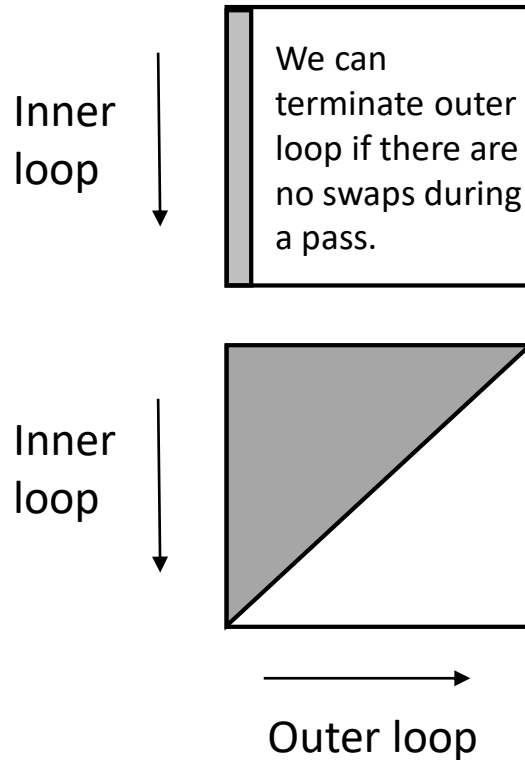
# Bubble Sort Algorithm

// You don't always need to make  $N - 1$  passes in outer loop.

```
for i = 0 to  $N - 2$  {  
    swapped = false  
    for k = 0 to  $N - 2 - i$  {  
        if ( list[k] > list[k+1] ) {  
            list.swap( k, k+1 )  
            swapped = true  
        }  
    }  
    if !(swapped)  
        break // return  
}
```

# Time Complexity ?

## Bubblesort



**Best case**

**Worst case**

Gray regions in the square indicate the indices examined through each pass through the inner loop.

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Lecture 12

Algorithms for Sorting a List:

bubble sort

selection sort

insertion sort

Feb. 2, 2022

# Selection Sort

Partition the list into two parts:

- the first part contains the smallest elements and is sorted
- the second part contains “the rest” of the elements  
(not necessarily sorted)

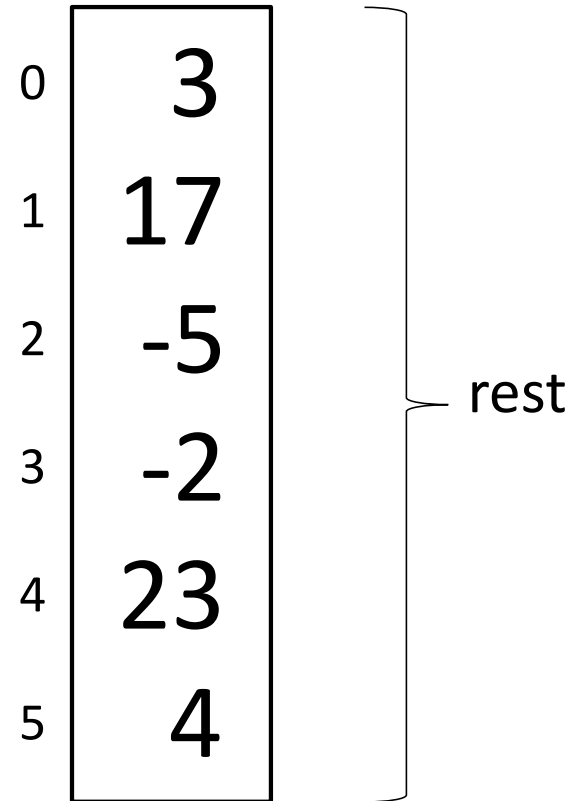
The sorted part is initially empty.

Repeat  $N - 1$  times {

- find the smallest element in “the rest”
- swap that element with the first element in “the rest”,
- this expands the first part of the list by 1

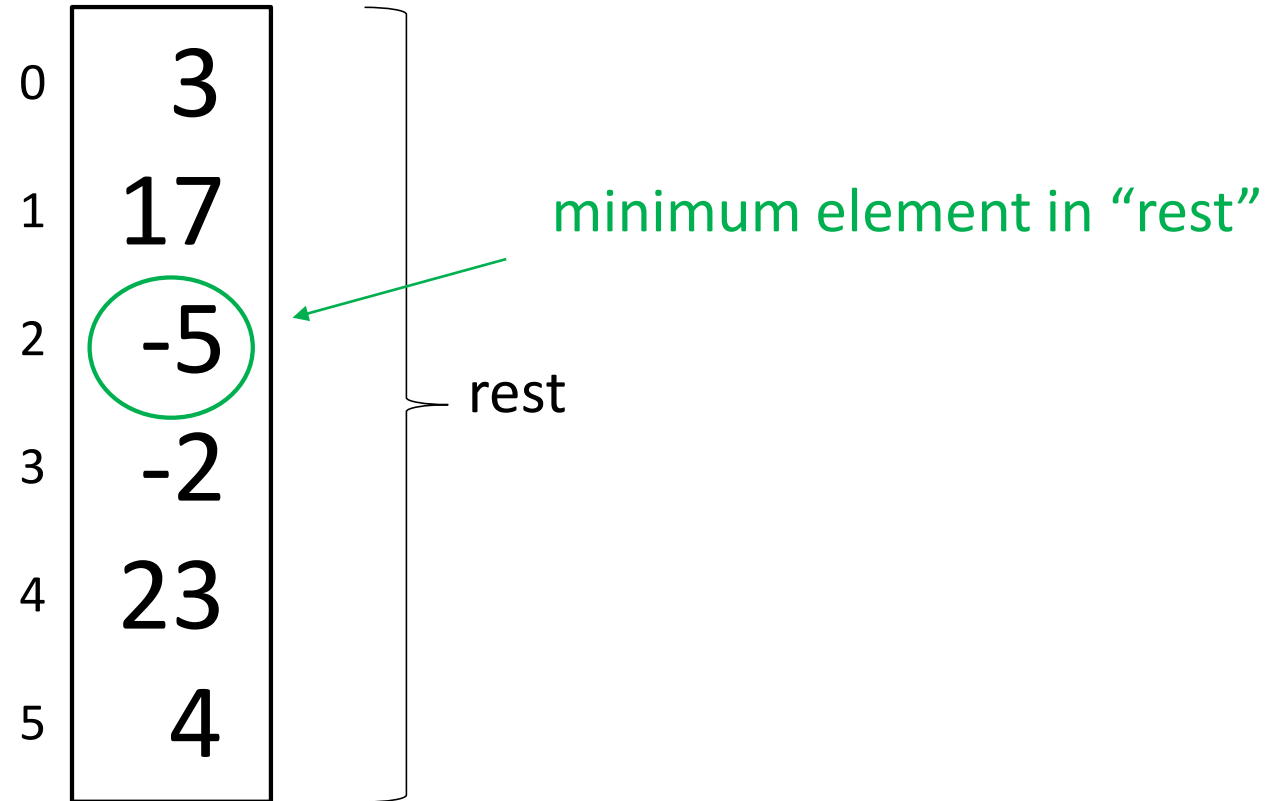
# Example

sorted part is empty



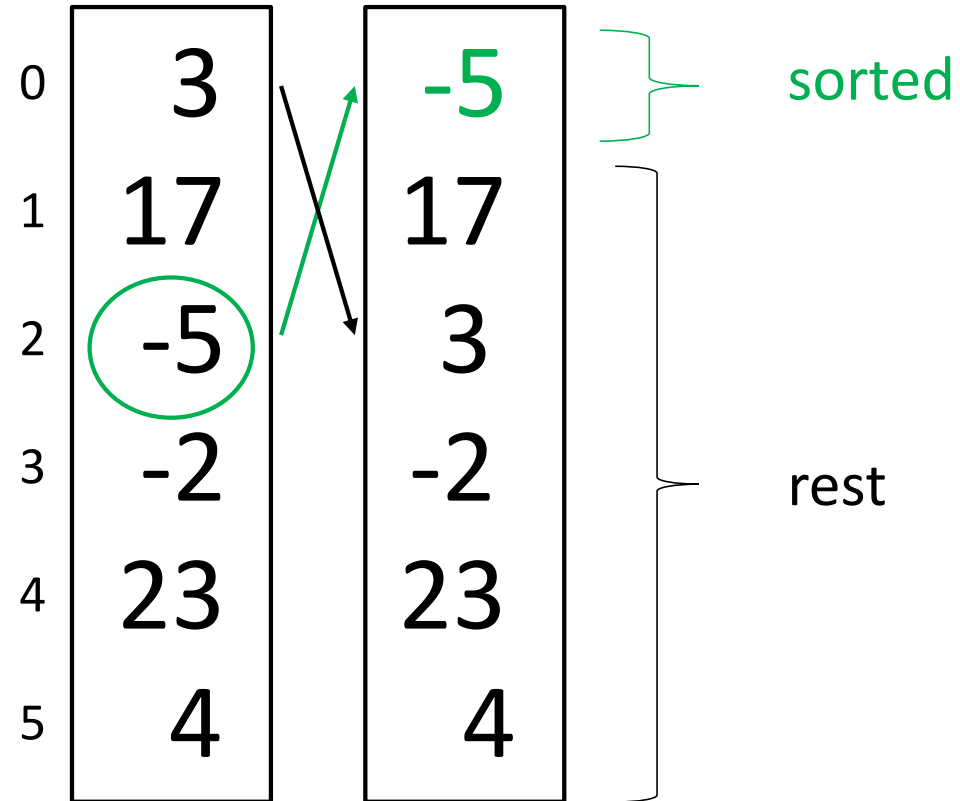
# Example

sorted part is empty



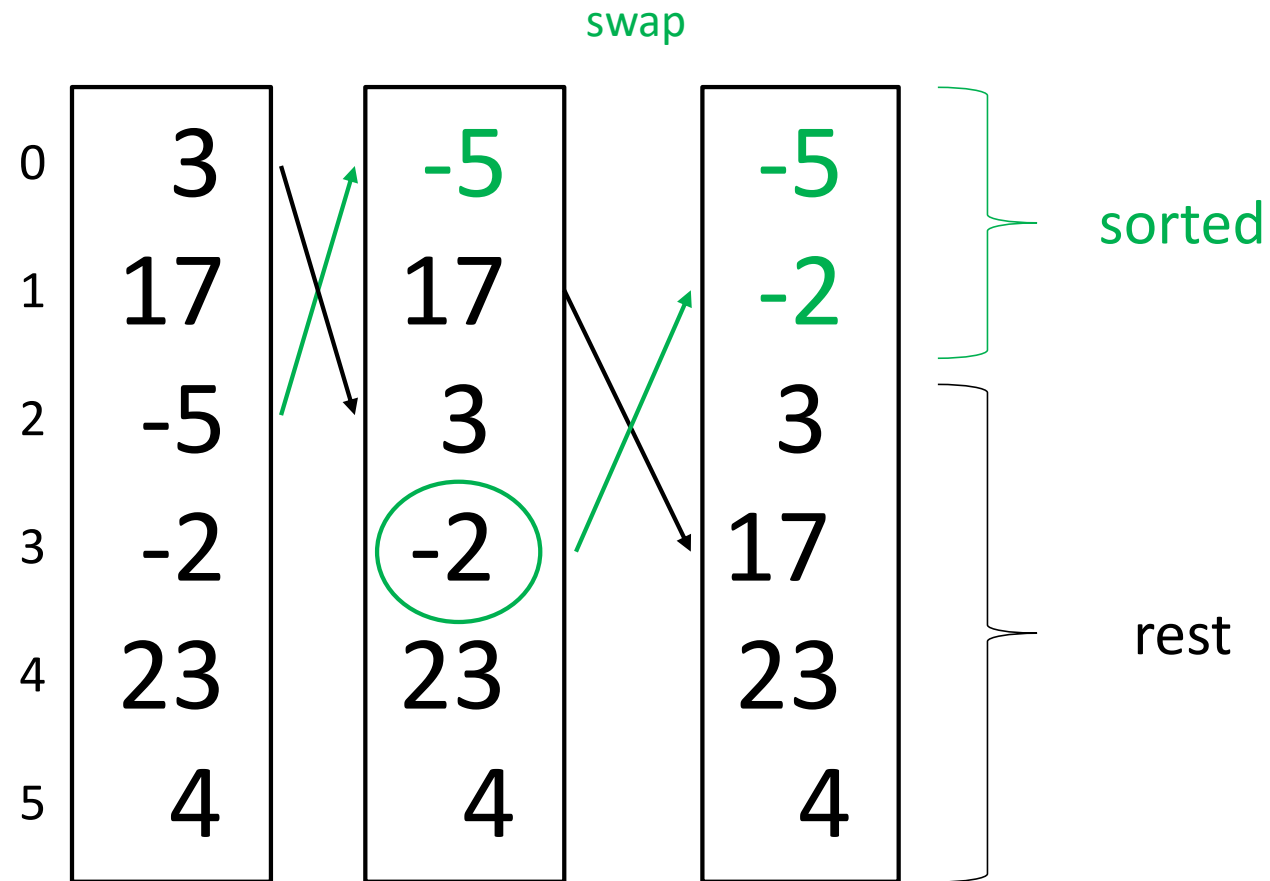
# Example

swap

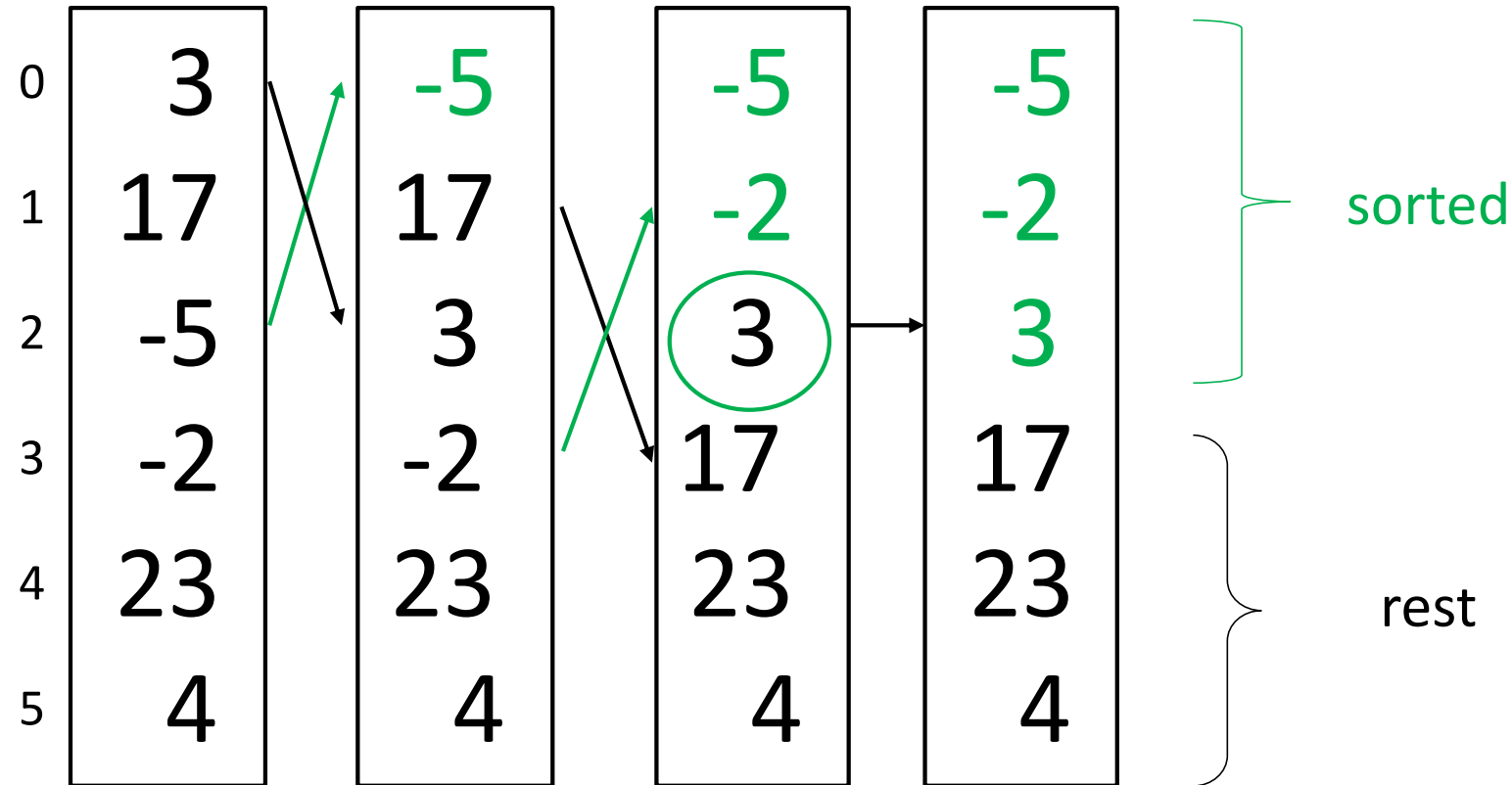




# Example

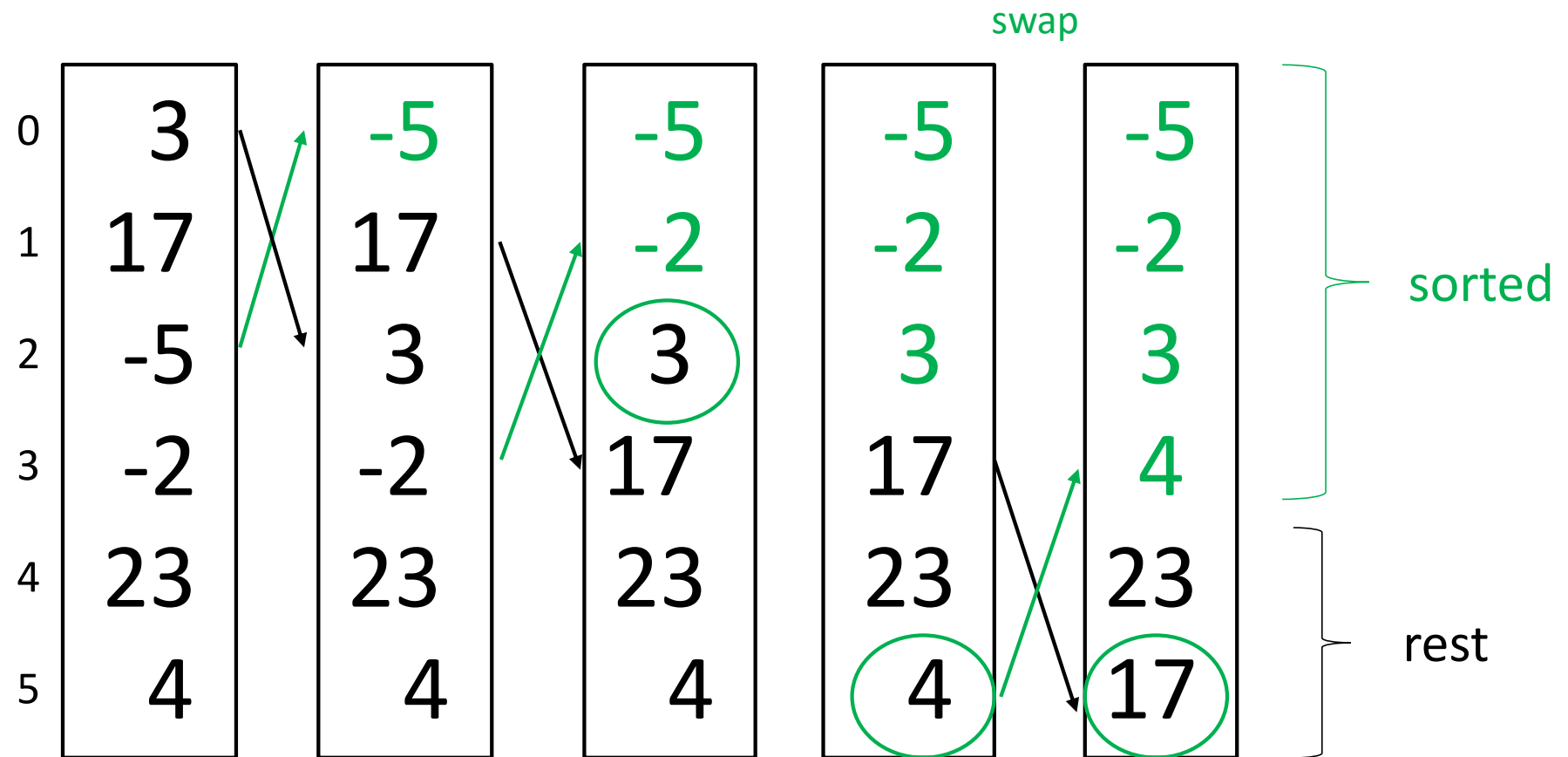


# Example

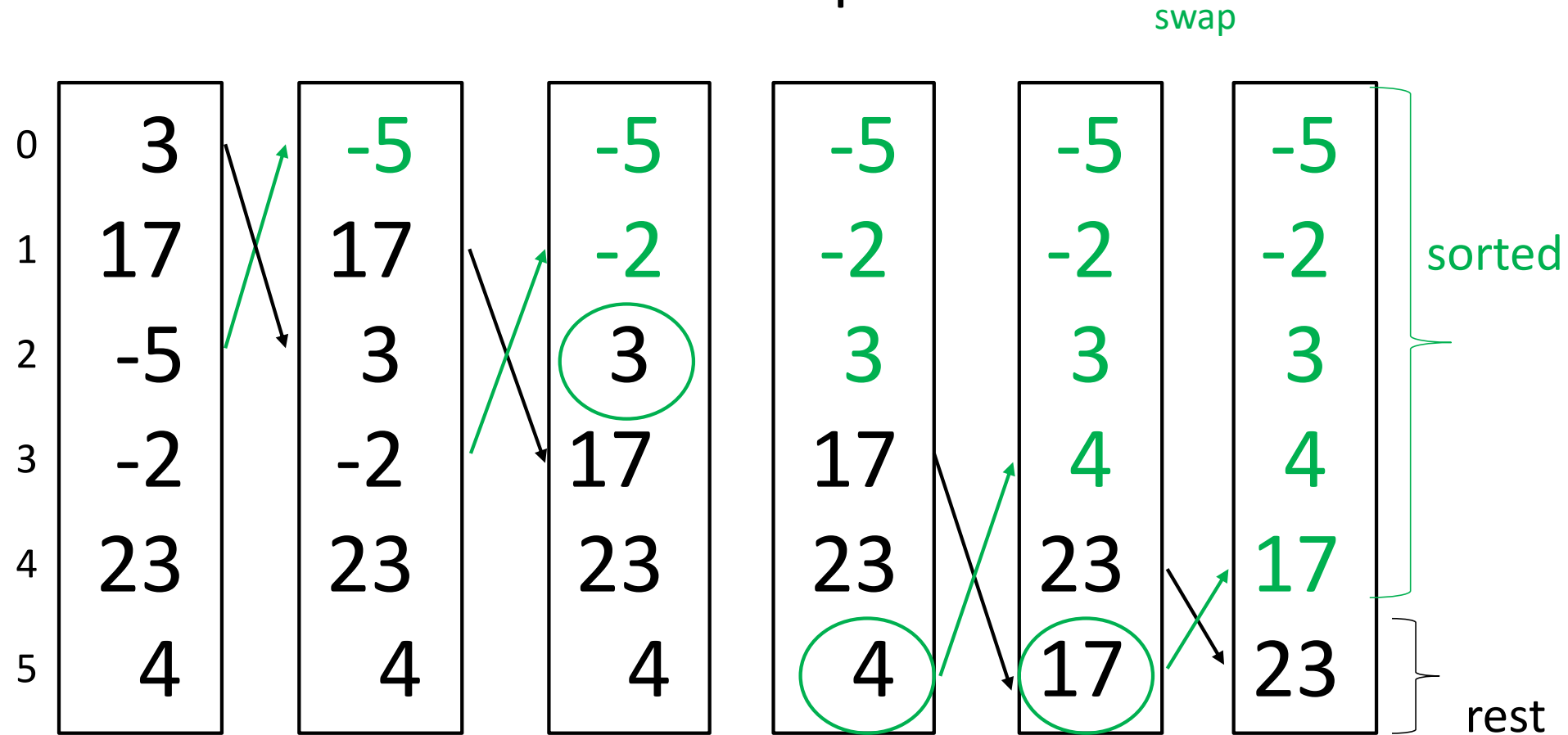


3 is the minimum  
element already

# Example



# Example



# Selection Sort

```
for  $i = 0$  to  $N - 2$  {
```

```
    index =  $i$ 
```

```
    minValue = list[  $i$  ]
```

```
    for  $k = i + 1$  to  $N - 1$  {
```

```
        if ( list[ $k$ ] < minValue ) {
```

```
            minValue = list[ $k$ ]
```

```
            index =  $k$ 
```

```
        }
```

```
    }
```

```
    if (  $i \neq$  index )
```

```
        list.swap( $i$ , index )
```

```
}
```

repeat  $N - 1$  times

Take the first element in the rest and let it be the temporary min value.

For each other element in rest,

if element is smaller than the min value, then it will become the new min value. So remember its index.

Swap (if it is necessary)

# Selection Sort

```
for  $i = 0$  to  $N - 2$  {
```

```
    index =  $i$ 
```

```
    minValue = list[  $i$  ]
```

```
    for  $k = i + 1$  to  $N - 1$  {
```

```
        if ( list[ $k$ ] < minValue ) {
```

```
            minValue = list[ $k$ ]
```

```
            index =  $k$ 
```

```
        }
```

```
    }
```

```
    if (  $i \neq$  index )
```

```
        list.swap( $i$ , index )
```

```
}
```



This is the bottleneck  
(the inner loop).

# Selection Sort

```
for  $i = 0$  to  $N - 2$   
    for  $k = i + 1$  to  $N - 1$   
        { .... }
```

Q: how many times does { .... } get executed?

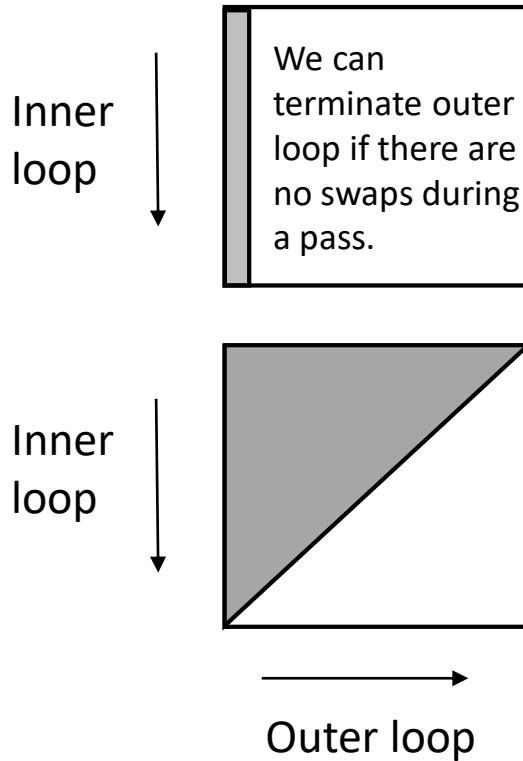
A:  $N - 1 + N - 2 + N - 3 + \dots + 2 + 1$

$\uparrow$              $\uparrow$              $\uparrow$                                      $\uparrow$

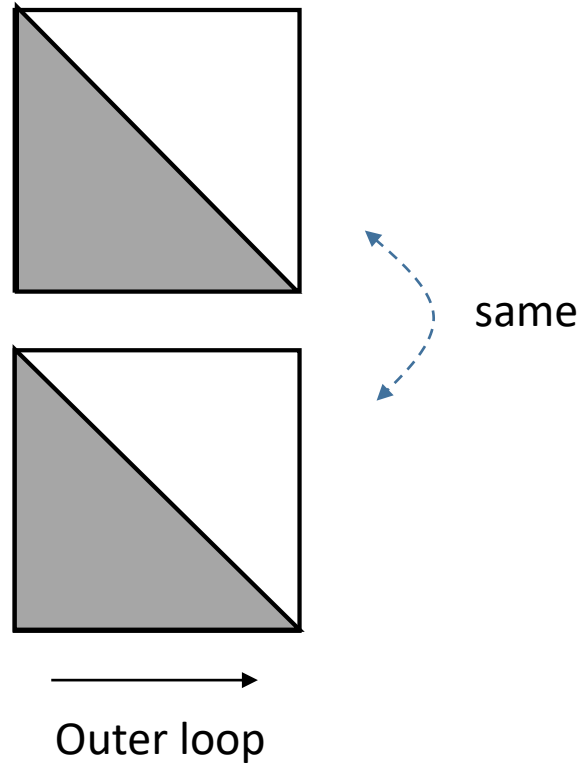
$i = 0$              $i = 1$              $i = 2$                                      $i = N - 2$

$$= N(N - 1) / 2$$

## Bubblesort



## Selection sort





# Insertion Sort

for  $i = 1$  to  $N - 1$  {

Insert element  $\text{list}[i]$  into its correct position with respect to the list elements at indices 0 to  $i - 1$

(At the start of pass  $i$ , the elements at indices 0 to  $i - 1$  are sorted *only amongst themselves*.

This is a weaker condition than in selection sort.)

}

Initial list

0	3
1	17
2	-5
3	-2
4	23
5	4

Initial list

$i = 1$

$i = 2$

$i = 3$

0	3	3	-5	-5
1	17	17	3	3
2	-5	-5	17	17
3	-2	-2	-2	-2
4	23	23	23	23
5	4	4	4	4

(At the start of pass 3, the elements at indices 0 to 2 are sorted *amongst themselves*.)

Initial list

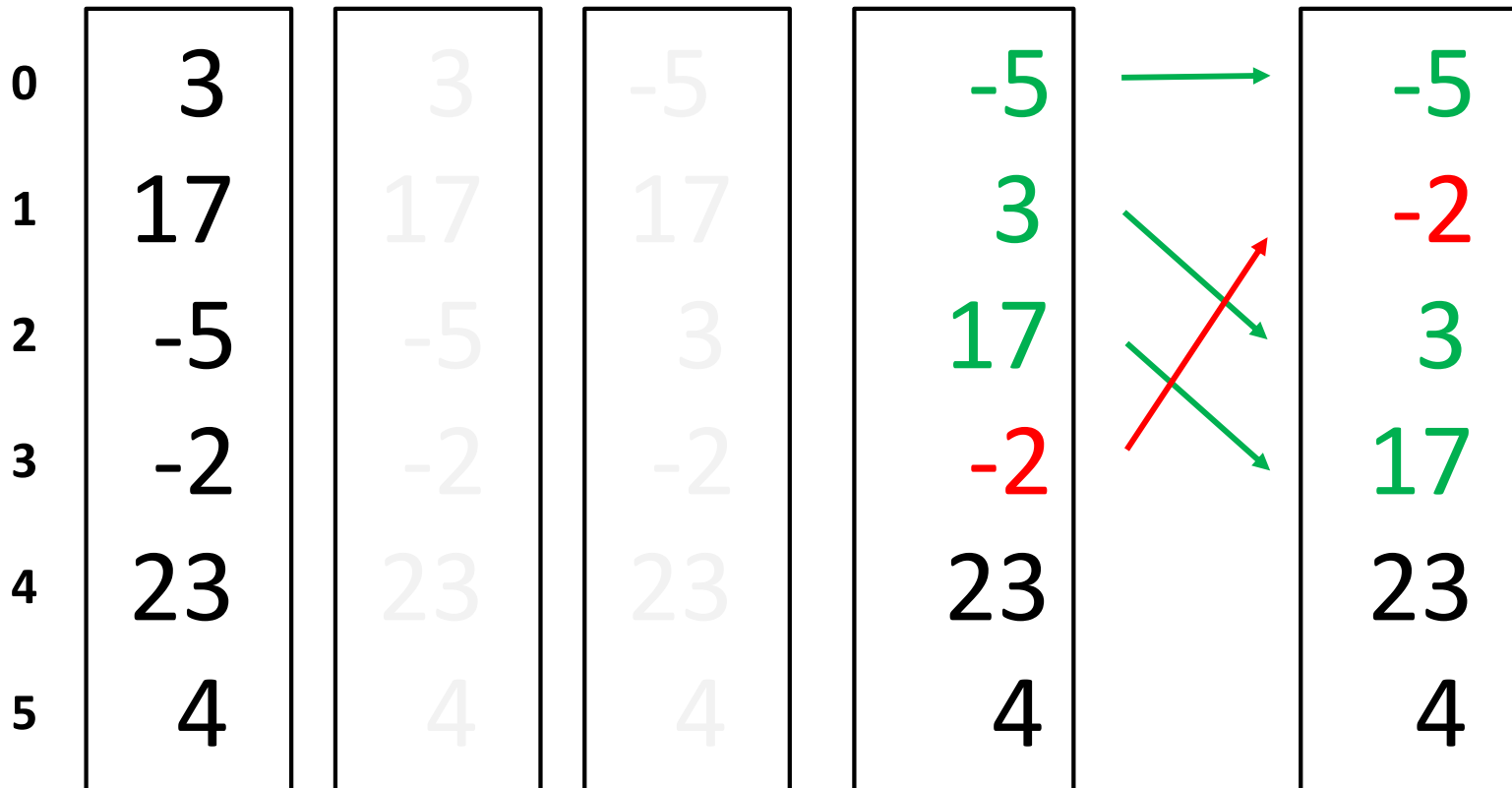
Insert element  $i$  into  
its correct position  
with respect to 0 to  $i-1$

$i = 1$

$i = 2$

$i = 3$

$i = 4$



(At the start of pass 4,  
the  
elements at  
indices 0 to  
3 are sorted  
*amongst  
themselves.*

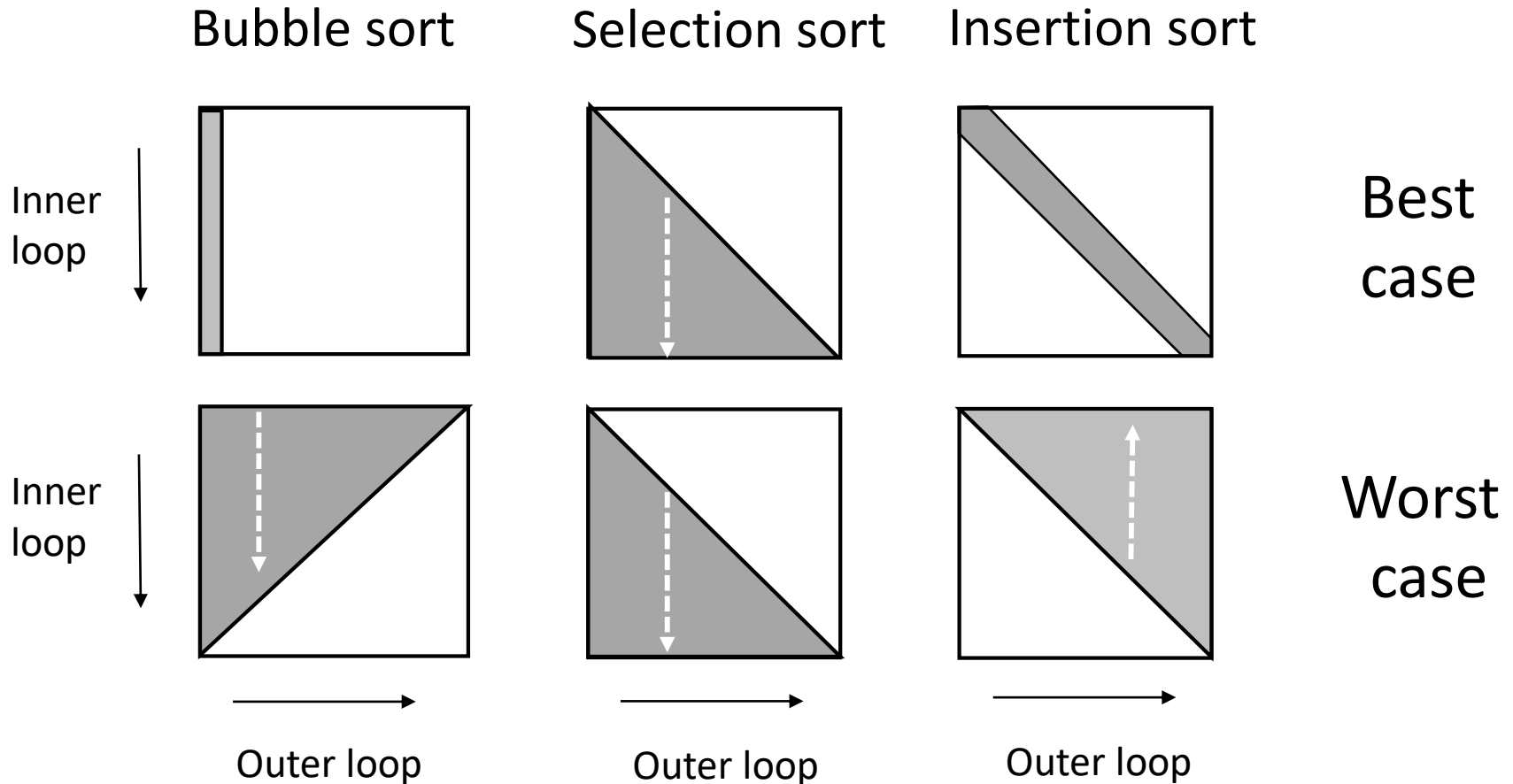
Mechanism is similar to inserting (adding) an element to an array list:

Shift all elements *forward* by one position to make a hole, and then fill the hole.

# Insertion Sort

```
for  $i = 1$  to  $N - 1$  { // index of element to move
    e = list[  $i$  ] // store as tmp
     $k = i$ 
    while ( $k > 0$ ) and (e < list[  $k - 1$  ]){
        list[ $k$ ] = list[ $k - 1$ ] // move it forward
         $k = k - 1$ 
    }
    list[ $k$ ] = e
}
```

# Time Complexity



Best case(s) : bubble and insertion sort are  $O(N)$ , selection sort is  $O(N^2)$ .  
Worst case : each of the three algorithms is  $O(N^2)$ .

# Sorting Algorithms

- Bubble sort
- Selection sort
- Insertion sort



today  $O(N^2)$

- Mergesort
- Heapsort
- Quicksort



lectures 22, 28  $O(N \log N)$



# Hector Tutorial TODAY on Zoom at 6 pm

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Assignments - A1 Michael Langer STAFF 3d

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A1 - Printing Course and usual NullPointe...  
Assignments - A1 Héctor Leos STAFF 2d

## A1 Tutorial #230



Héctor Leos STAFF

a day ago in Assignments - A1



UNPIN



STAR



Hey all!

I'll host a tutorial for A1, mainly to help you conceptually, go over some examples, and take some advice about how to implement your solutions. The tutorial will take place **tomorrow, Feb 2, at 6pm**. Here's the link to join: <https://mcgill.zoom.us/j/6327362007>