

Lecture 6

- Conclusion of KR
- Search
 - Blind search

Concept Description Language

- A specialized language for efficient inference.
- Represent
 - classes of objects,
 - sub-classes of classes,
 - instances of classes,
 - properties of instances (and classes).
- Akin to inheritance in object-oriented programming.
- A **semantic network** is a graph-based representation that addresses the same idea.

(See DAA pp. 107-109.)

Nonmonotonic Logic

A monotonic logic:

Things that are theorems remain theorems as we add additional axioms.

Non-monotonic logic: formulas that were once theorems may not remain so as the theory is augmented.

- Idea: add “default” assumptions to the theory in the absence of complete knowledge.
- These assumptions may be **retracted** later, as we learn more.
- As a result, conclusions based on those assumptions are invalidated.

Minimal (nonmonotonic) models

- Can induce a preference ordering on interpretations.

Deductive retrieval

- Deductive retrieval uses a KB to store information, and uses rules to achieve goals.
- A system for maintaining a knowledge base. Includes retracting conclusions based on information that changes or which is deleted.
- Use **forward chaining** and **backward chaining**.
 - Forward: start from KB and see what can be inferred (leading to the goal, we hope). Especially happens when a question involves new knowledge.
 - Q. if pigs fly then will I pass?
 - Backwards: we already saw **goal reduction**.
 - Start from the goal.
 - See what we need to infer it.

Search

- Reference: DAA Chapter 4.
- The process of explicitly examining a set of objects to find a particular one, or satisfy some goal.
- In the context, the objects are typically possible configurations of some problem representation.
- Eg.
 - Towers of Hanoi problem.
 - Path planning
 - Theorem proving!
 - Game playing.

Search

- Search is a central topic in AI
 - Originated with Newell and Simon's work on problem solving.
Famous book:
``Human Problem Solving" (1972)
- Automated reasoning is a natural search task
More recently: Given that almost all AI formalisms (planning, learning, etc.) are NP-Complete or worse, some form of search is generally **unavoidable** (no ``smarter" alg. Available)

State space

- In AI, search usually refers to search of a state space.
- State space: the ensemble of possible configurations of the domain of interest.
 - Like phase space in physics
- E.g.
 - Chess: The set of allowed arrangements of pieces on a chess board.
 - Speech understanding: The set of possible arrangements of words that make valid sentences.
 - Robot motion planning: the set of accessible & safe locations for the robot.

Search: Problem Definition

State space -- described by an **initial state** and the set of possible actions available (**operators**).

- A path is any sequence of actions that lead from one state to another.

Goal test -- applicable to a single state to determine if it is a (the) goal state.

Path cost -- relevant if more than one path leads to the goal, and we want the shortest path.

- Note: very general formulation. Can be somewhat unnatural.

Graphs, Trees & Search

- We can visualize generic state space search in terms of searching a graph or tree.
- Graph search corresponds to looking for a particular state given an arbitrary transition diagram.
 - A graph is defined as $G = (V, E)$
 - V : set of vertices (i.e. states)
 - E : set of transitions $e_i = (v_j, v_k)$. Can be directed or undirected.
- Tree search results when there is a natural ordering, as with game playing.

Traversal

- To traverse means to visit the vertices in some systematic order. You should be familiar with various traversal methods for trees:

preorder: visit each node before its children.

postorder: visit each node after its children.

inorder (for binary trees only): visit left subtree, node, right subtree.

Familiar ideas....

- **Blind search:** just examine “successive” alternative possibilities.
- Does not exploit knowledge of what states to examine first.
 - It what order should we consider the states?
- Sequentially along a path? Leads to **Depth First Search** (DFS)
- A little bit along each possible path in turn? Leads to **Breadth First Search** (BFS).

Depth First Search

- Key idea: pursue a sequence of successive states *as long as possible*.

```
unmark all vertices
```

```
  choose some starting vertex x
```

```
  mark x
```

```
  list L = x
```

```
  tree T = x
```

```
  while L nonempty
```

```
    choose the vertex v from front of list
```

```
    visit v
```

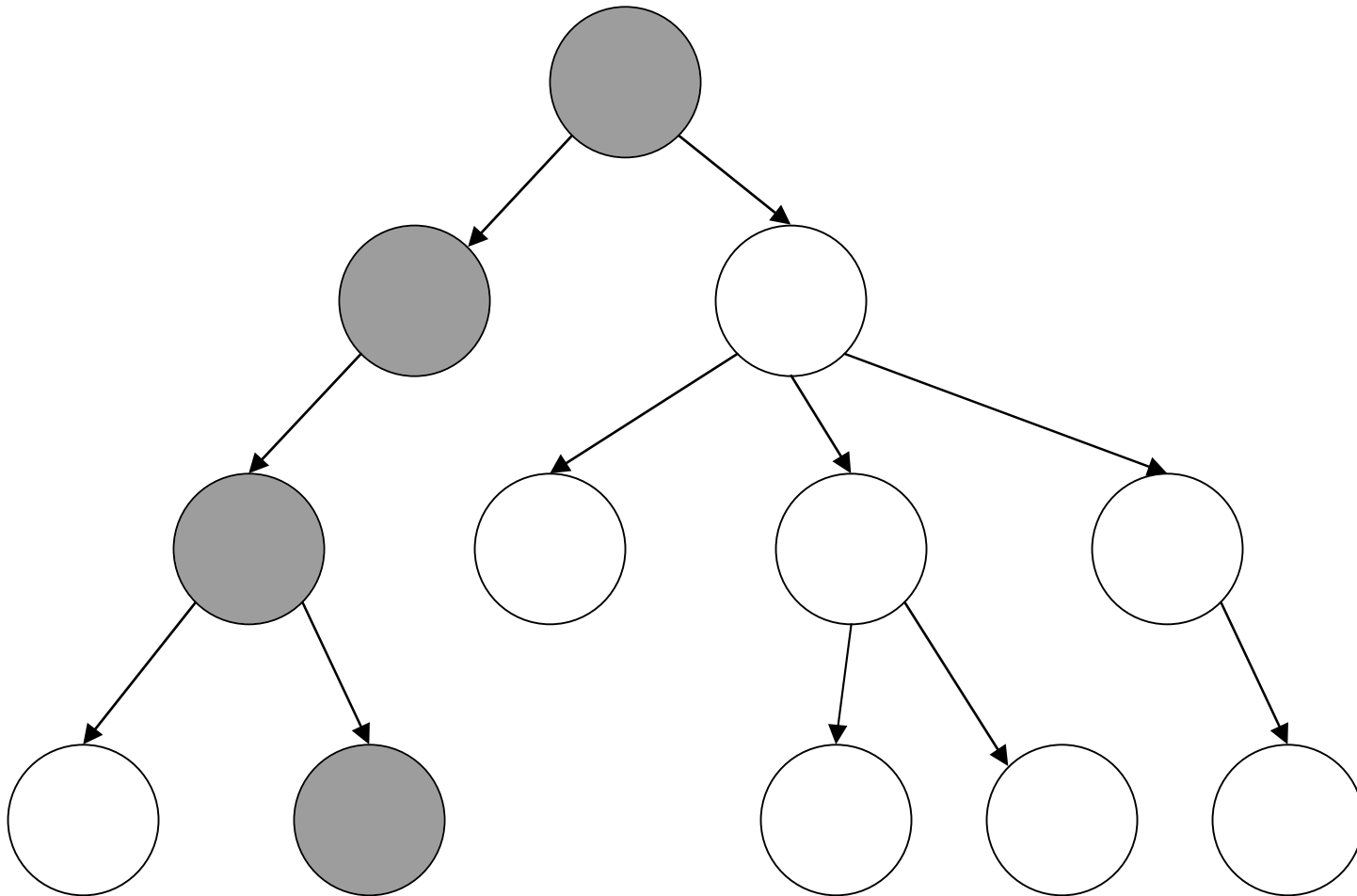
```
    for each unmarked neighbor w
```

```
      mark w
```

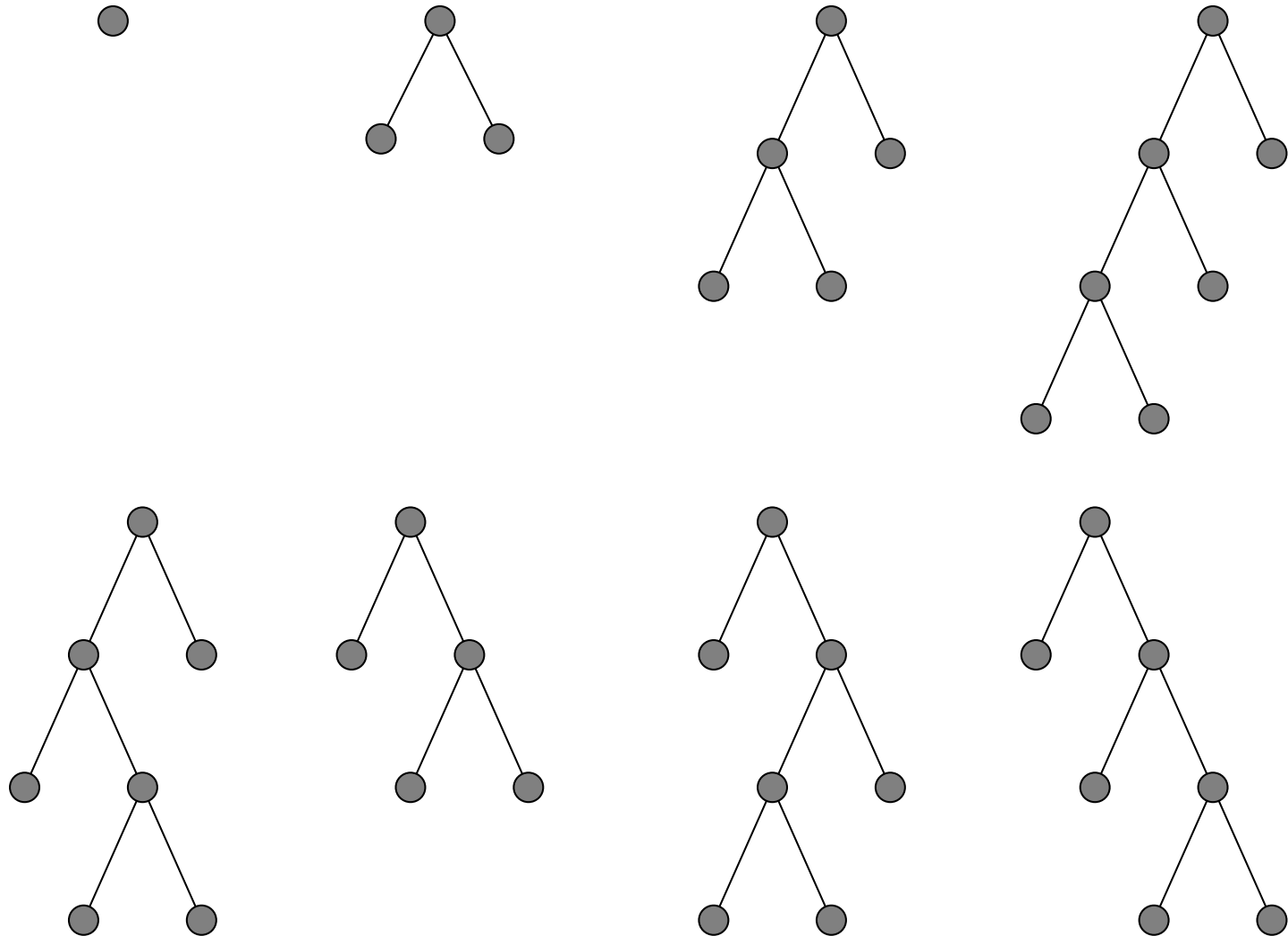
```
      add it to front of list
```

```
      add edge (v,w) to T
```

Depth First Search



DFS illustrated



BFS

- Key: explore nodes at the same distance from the start at the same time

```
unmark all vertices
```

```
  choose some starting vertex x
```

```
  mark x
```

```
  list L = x
```

```
  tree T = x
```

```
  while L nonempty
```

```
    choose the vertex v from front of list
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```
    visit v
```

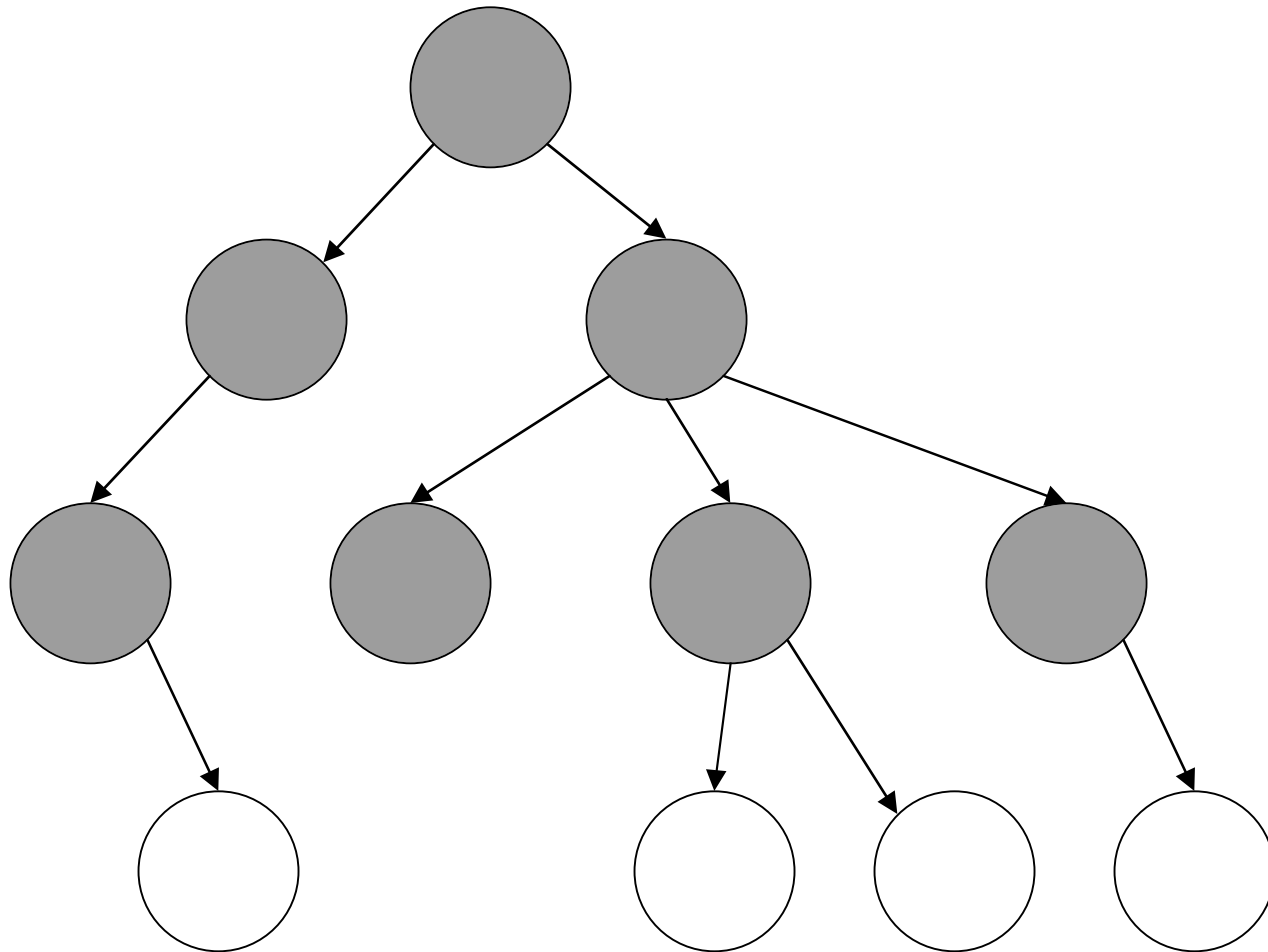
```
    for each unmarked neighbor w
```

```
      mark w
```

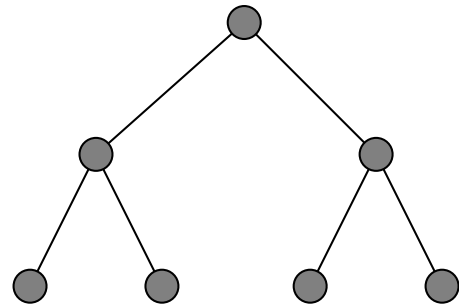
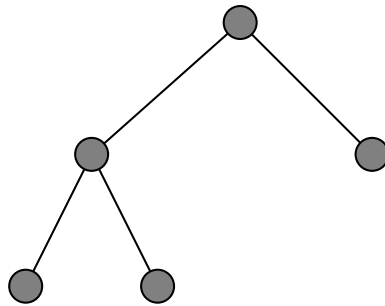
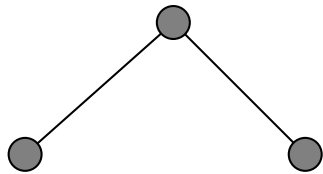
```
      add it to end of list
```

```
      add edge (v,w) to T
```


Breadth First Search



BFS illustrated



Key issues in search

- Here's what to keep in mind.
- Completeness: are we assured to find a solution (if one exists)?
- Space complexity: how much storage do we need?
- Time complexity: how many operations do we need?
- Solution quality: how good is the solution
- Also...
 - Can we expect an interim solution?
 - Can we refine a partial/inadequate solution?
 - Can we cope with imperfect knowledge?

Example I : Cryptarithmic

$$\begin{array}{r} \text{SEND} \\ + \text{MORE} \\ \hline \text{MONEY} \end{array}$$

- Find substitution of digits for letters such
- that the resulting sum is arithmetically correct.
- Each letter must stand for a different digit.

Cryptarithmic, cont.

- **States**: a (partial) assignment of digits to letters.
- **Operators**: the act of assigning digits to letters.
- **Goal test**: all letters have been assigned digits and sum is correct.
- **Path cost**: zero. All solutions are equally valid.

- Solution method?
- Depth first search:

Search performance

- Key issues that determine the nature of the problem are:
 - **Branching factor** of the search space: how many options do we have at any time?
 - Typically summarized by the worst-case branching factor, which may be quite pessimistic.
 - **Solution depth**: how long is the path to the first solution?

Example: knight's tour

- Tour executed by a chess knight to cover (touch) every square on a chess board.
 - Sub-problem: Find a path from one position to another.
- States: possible positions on the chess board.
- Operators: the ways a knight moves.
- Goal: the positions of the pawns.
- Path cost: the number of moves used to get to a position.

Knight's

- How large is the state space?
 - A knight has up to 8 moves per turn.
 - Each possible tour must be verified up to the end of the trip. For a board of width N , there are $N*N$ squares.
 - Thus, each tour can be up to $N*N$ states in length.
 - If the correct solution is found last, we might consider up every wrong tour first.
- $O(8^{N^2})$ states to examine!

What does this say about BFS? DFS?

Example: Knight's heuristics

- **Zero:** a board can be considered a dead end if any square has zero paths remaining to it. Any square with no paths to it would be unreachable, so no Knight's Tour would exist.
- **Two ones:** a board with more than one square with only one path to it is a dead end. A square with one path left is necessarily a dead end, so two of them indicate a dead end position.
- **Move to one:** the move finder should never choose a move to a square with one path left unless it is the last move; such a choice would otherwise lead to a dead end.

- How well will we do if we use “blind” DFS?
- That is, if we consider random tours?

Applet

Full tour produced at 3724122

The Distribution of Tour Length:

1	2	3	4	5	6	7	8
0	0	0	2184	2048	6187	5890	9887
9	10	11	12	13	14	15	16
10099	14776	14734	20561	19899	27129	26337	34769
17	18	19	20	21	22	23	24
34021	44222	42591	54671	52573	65745	63380	78348
25	26	27	28	29	30	31	32
74752	91527	87630	105141	100167	119356	112681	132464
33	34	35	36	37	38	39	40
124064	143282	132906	150826	137549	153940	138422	151938
41	42	43	44	45	46	47	48
133214	142374	121666	126676	104952	104410	83146	79196
49	50	51	52	53	54	55	56
59879	53945	37838	31949	20887	15887	9171	6208
57	58	59	60	61	62	63	64
3063	1784	723	322	96	36	2	1

Total elapsed time: 4.8029885 hour

- Three heuristics based on the number of paths remaining to each square were implemented and tested in combination, as well as a representational speedup. The optimal combination of the heuristics is to eliminate boards with either a square with zero paths remaining or a square with two ones remaining; this combination led to a 950-fold speedup on a 6x6 board. The representational speedup led to a 2.5-fold additional speedup on a 6x6 board.
 - Michael Bernstein

Heuristics: Knight's Performance

ZERO	TWO ONES	MOVE TO ONE	Time
0	0	0	9500 ms
1	0	0	415 ms
0	1	0	13.7 ms
0	0	1	6430 ms
0	1	1	13.2 ms
1	1	0	10.0 ms
1	0	1	350 ms
1	1	1	10.4 ms

BFS

- Consider a state space with a uniform branching factor of b
- At each level we have b nodes for every node we had before

$$1 + b + b^2 + b^3 + b^4 \dots + b^d$$

So, solution at depth d implies we must expand

$O(b^d)$ nodes!

- internal nodes $(b^{d+1}-1)/(b-1)$
- Leaf nodes: $b^{d+1}/2$
- why 2?

Depth	Nodes	Time	Memory
0	1	1 millisecond	100 bytes
2	111	.1 seconds	11 kilobytes
4	11,111	11 seconds	1 megabyte
6	10^6	18 minutes	111 megabytes
8	10^8	31 hours	11 gigabytes
10	10^{10}	128 days	1 terabyte
12	10^{12}	35 years	111 terabytes
14	10^{14}	3500 years	11,111 terabytes