# Lecture 5

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## Today's lecture

- Administrative issues
  - Comments on assignment
  - PDF files
  - Class notes
- Knowledge representation: wrap-up
  - Prolog details
  - Non-monotonic logic
  - Forward and backward chaining
- Introduction to search



## Don't care

- Symbol _ (underscore) is used to match a predicate that we don't plan to use on the right-hand side.
- It's like a dummy variable.

Eg. **likes(a,b).**
Would return **true** no matter what a & b are.
We can use **likes(a,_).**
…..or **likes(_,_).**



## Prolog (continued)

- Supports **lists** of items
  - [] - empty list
  - [1,2,3] - 3 items
  - [bob, ted, alice] - three objects
  - [[a], [1,2,3], []] - a list of lists

To examine a sub-part
[ H | L ]
refers to a list decomposed into a
head:H (the first element)
and a remaining part
tail:T

## Lists: [H|T]

[1,2,3]
- head is **1**

[bob, ted, alice]
- head is **bob**

[ [a], [1,2,3], [ ] ]
- head is **[a]**

[ [ [1,2,3]], [1,2,3], [ ] ]
- head is [ **[1,2,3] ]**

[ ]
- cannot match



## Membership: the body.

- Step 2: if it's not the head, then there must be a sublist for which it is the head.
  - Recursive definition
- See if the item is the head of the tail portion.

```
member(Item,[Head|Tail]) :-
  member(Item,Tail).
```



## Testing membership

- Now we can easily define a predicate to test for list membership.
- Step 1: the head

  **member(H,[H|L]).**
  - First argument is an item.
  - Second argument is a list.
    - This matches if H is the head of the list.

  member(bob,[bob,alice]) unifies with member(H,[H|L)) is we let bob match H and [bob,alice] match [H|L].

  **member( H, [H | _ ] ).**



## Membership: complete

```
member( Item, [ Item | _ ] ).
member(Item, [ _ | Rest ] ) :-
   member(Item, Rest).
```

## Unification: examining combinations

- **Remember: prolog execution proceeds by repeated unifications, applied recursively.**
- **Consider:**
  - **foo(x) :- bar(x).**
  - **bar(x) :- foo(x).**
  - **This will lead to a problem: the unification will never terminate!**



## Improved foo!

```
member…

foo(X,L) :-  not(member(X,L)),
      bar(X,[X|L]).

bar(X,L) :-  not(member(X,L)),
      foo(X,[X|L]).

foo(a).
```

Improved foo!

foo(y,[]).



## Recursion fix

- How can we fix the infinite recursion?
  - Never re-examine an already-considered unifier (i.e. solution).

1. Within the definition, save the previous solutions (unifications).
2. Check if the new unifier (solution) is one of those.

How?

Use a list!

```
foo(X,L) :- …
```

X is the item,

L is a list of prior unifiers



## Concept Description Language

- A specialized language for efficient inference.
- Represent
  - classes of objects,
  - sub-classes of classes,
  - instances of classes,
  - properties of instances (and classes).
- Akin to inheritance in object-oriented programming.

- A **semantic network** is a graph-based representation that addresses the same idea.

(See DAA pp. 107-109.)