

## 1.3 Definitions for “machine,” “mechanism,” and “linkage”:

**Machine**

- Definitions in Merriam Webster’s Collegiate Dictionary (on-line, 2002):
  - (archaic): a constructed thing whether material or immaterial.
  - an assemblage of parts that transmit forces, motion, and energy one to another in a predetermined manner
  - an instrument (as a lever) designed to transmit or modify the application of power, force, or motion
  - a mechanically, electrically, or electronically operated device for performing a task (a calculating machine, a card-sorting machine)

*Comment: comprehensive definitions when considered as a whole*

- An *apparatus* for transformation of power, materials, and information to substitute or simplify physical or intellectual work (Frolov, 1987). *Comment: a comprehensive definition, that includes computers*
- *Mechanical system* that performs a specific task, such as the forming of material, and the transference and transformation of *motion* and *force*, Vol. 38, Nos. 7–10 (2003) of *Mechanism and Machine Theory* on Standardization of Terminology. *Comment: leaves computers out*
- An *apparatus* for applying mechanical power, having several parts, each with definite function (The Concise Oxford Dictionary). *Comment: same as above*
- An *apparatus* consisting of interrelated parts with separate functions, used in the performance of some kind of work (The Random House College Dictionary). *Comment: ditto*
- Any *system* in which a specific correspondence exists between an input form of energy or information and the corresponding ones at the output (Loosely translated from Le Petit Robert). *Comment: as comprehensive as Frolov’s*

**Mechanism**

- A piece of machinery (Merriam Webster’s Collegiate Dictionary (on-line, 2002)). *Comment: too vague*
- Definitions in Vol. 38, Nos. 7–10 (2003) of *Mechanism and Machine Theory* on Standardization of Terminology.
  - *System* of bodies designed to convert *motions* of, and *forces* on, one or several bodies into constrained motions of, and *forces* on, other bodies. *Comment: English could be terser, but idea is fine.*
  - *Kinematic chain* with one of its components (*link* or *joint*) connected to the *frame*. *Comment: confuses mechanism with its representation as a kinematic chain*
- Structure, adaptation of parts of machine; system of mutually adapted parts working together (as) in machine (The Concise Oxford Dictionary)
- An assembly of moving parts performing a complete functional motion (The Random House College Dictionary)
- A combination layout of pieces or elements, assembled with the goal of (producing) an operation as a unit (Loosely translated from Le Petit Robert)

*Comment: In all above definitions, the concept of goal or task is present*

**Linkage**

- Definitions in Merriam Webster’s Collegiate Dictionary (on-line, 2002):
  - a system of links. *Comment: concise and comprehensive*
  - a system of links or bars which are jointed together and more or less constrained by having a link or links fixed and by means of which straight or nearly straight lines or other point paths may be traced. *Comment: unnecessarily cumbersome and limited to path-generating linkages*
- *Kinematic chain* whose *joints* are equivalent to *lower pairs* only (Vol. 38, Nos. 7–10 (2003) of *Mechanism and Machine Theory* on Standardization of Terminology). *Comment: confuses linkage with its representation*

Frolov, K. V., (editor), 1987, *Teoriya Mechanismov i Mashin* (Theory of Mechanisms and Machines), Vyschaya Shkola, Moscow (in Russian)

*Mechanism and Machine Theory* on Standardization of Terminology, 2003, Vol. 38, Nos. 7–10

*Le Petit Robert*, 1994, Dictionnaires Le Robert, Paris

*Random House Webster's College Dictionary*, 1997, Random House, New York

*Merriam Webster's Collegiate Dictionary*, on-line 2002

**1.10** Here we want to estimate the time required to add two floating-point numbers. Since this time is very small, it is more suitable to perform a bunch of additions, say  $10^7$ , which should amount to a total time of the order of seconds.

- The Microsoft Visual C++6.0 program below was run on a Pentium IV 2.0:

```
#include <iostream>
#include <time.h>
int main()
{
    clock_t start, end;
    start = clock();
    float a, i, CLOCKS_PER_mSEC=CLOCKS_PER_SEC/1000;
    for(i=1; i<=10000000; i++)
        a = 5+5;
    std::cout << "a= " << a << std::endl;
    end = clock();
    long duration = (long)(end-start)/CLOCKS_PER_mSEC;
    std::cout << "Time: " << duration << "ms" << std::endl;
    return 0;
}
```

A Pentium IV 2.0 processor required 125 ms to perform the  $10^7$  additions, which amounts to  $1.25 \times 10^{-8}$  s/add, or  $8 \times 10^7$  add/s.

- The C program below was run for the same purpose on the CLUMEQ supercomputer, AMD Athlon 1900+ cluster:

```
#include "stdio.h"
#include "math.h"
#include <sys/time.h>
```

```

#include <sys/resource.h>
#define RUSAGE_SELF 0 /* calling process */
main()
{
    int getrusage(int who, struct rusage *rusage);
    long diffsec, diffmsec;
    float a;
    int i;
    struct rusage time_begin, time_end;
    struct rusage *pb, *pe;
    pb=&time_begin;
    getrusage(RUSAGE_SELF, pb);
    for(i=1; i<=10000000; i++)
    {
        a = 5+5;
    }

    printf("a=%f \t\n",a);
    pe=&time_end;
    getrusage(RUSAGE_SELF,pe);

    diffsec=time_end.ru_utime.tv_sec-time_begin.ru_utime.tv_sec;
    diffmsec=time_end.ru_utime.tv_usec-time_begin.ru_utime.tv_usec;
    printf("User time used:%d microsec\n",diffmsec);
}

```

The CLUMEQ supercomputer required 40 ms to perform the  $10^7$  additions, which amounts to  $4 \times 10^{-9}$  add/mult, or  $2.5 \times 10^8$  add/s.

- 1.11 What we must find here is the largest floating-point number  $\epsilon$  that, when added to a given number  $a$ , leaves this number unchanged, i.e.,

$$a + \epsilon = a$$

- The Microsoft Visual C++6.0 program below was run on a Pentium IV 2.0 processor at 550 MHz:

```

#include <iostream>
#include <time.h>
int main()
{
    long i;
    //double a=1;
    float a=1;
    for(i=1; i<=10000000; i++)
    {
        a = a/10;
        if (a==0)
            break;
    }

    std::cout << "smallest floating-point"<< i-1 << std::endl;
    return 0;
}

```

The value of  $\epsilon$  reported in single-precision arithmetic was  $1.0e - 045$ ; in double-precision.....  
 $1.0e - 323$ .

- The program below was run on the CLUMEQ supercomputer, AMD Athlon 1900+ cluster:

```
#include "stdio.h"
#include "math.h"
#include <sys/resource.h>
main()
{
    int i;
    float a=1;
    for(i=1; i<=10000000; i++)
    {
        a=a/10;
        if (a==0)
            break;
    }

    printf("smallest floating-point=%d \n",i-1);
}
```

This machine reported exactly the same values as its Pentium counterpart.

**2.7** The determinant of a matrix is equal to the product of its eigenvalues. Thus

$$\det(\mathbf{1} + \mathbf{u}\mathbf{v}^T) = \det(\mathbf{T}) = \lambda_1 \lambda_2 \lambda_3 = 1 + \mathbf{v}^T \mathbf{u} = 1 + \mathbf{u} \cdot \mathbf{v}$$

**2.14** We have, from Cayley's Theorem,

$$\mathbf{Q} = (\mathbf{1} - \mathbf{S})(\mathbf{1} + \mathbf{S})^{-1}$$

where  $\mathbf{1} + \mathbf{S}$  is invertible. Indeed, if we write  $\text{vect}(\mathbf{S}) = \mathbf{s} \equiv \sigma \mathbf{e}$ , with  $\sigma > 0$  denoting the Euclidean norm of  $\mathbf{s}$ , then  $\mathbf{S} = \sigma \mathbf{E}$ , in which  $\mathbf{E}$  is the cross-product matrix of the unit vector  $\mathbf{e}$ . Hence, the eigenvalues of  $\mathbf{S}$  are, from Exercise 2.10,  $0$ ,  $\sigma j$ , and  $-\sigma j$ . Consequently, the eigenvalues of  $\mathbf{1} + \mathbf{S}$  are  $1$ ,  $1 + \sigma j$ , and  $1 - \sigma j$ , none of which is zero, thereby showing that  $\mathbf{1} + \mathbf{S}$  is indeed invertible. Multiplying both sides of this equation by  $(\mathbf{1} + \mathbf{S})$ , from the right, we obtain

$$\mathbf{Q}(\mathbf{1} + \mathbf{S}) = (\mathbf{1} - \mathbf{S})$$

i.e.,

$$(\mathbf{1} + \mathbf{Q})\mathbf{S} = \mathbf{1} - \mathbf{Q}$$

and hence,

$$\mathbf{S} = (\mathbf{1} + \mathbf{Q})^{-1}(\mathbf{1} - \mathbf{Q}).$$

For this factoring to be possible, then, matrix  $\mathbf{1} + \mathbf{Q}$  should be invertible, i.e., nonsingular. Therefore, the above-mentioned factoring is not possible if one of the eigenvalues of  $\mathbf{Q}$  is  $-1$ , which happens iff the angle of rotation is  $\pi$ .

**2.10** Let  $\lambda$  be an eigenvalue of  $\mathbf{E}$ , the cross-product matrix of the unit vector  $\mathbf{e}$ , and  $\mathbf{u}$  the corresponding unit eigenvector, i.e.,

$$\mathbf{E}\mathbf{u} = \lambda\mathbf{u} \quad (1)$$

Multiplying eq.(1) by  $\mathbf{E}$ , we obtain

$$\mathbf{E}^2\mathbf{u} = \lambda^2\mathbf{u} \quad (2)$$

and multiplying a second time by  $\mathbf{E}$ ,

$$\mathbf{E}^3\mathbf{u} = \lambda^3\mathbf{u} \quad (3)$$

However,

$$\mathbf{E}^2 = -\mathbf{1} + \mathbf{e}\mathbf{e}^T$$

and

$$\mathbf{E}^3 = \mathbf{E}(-\mathbf{1} + \mathbf{e}\mathbf{e}^T) = -\mathbf{E} + \underbrace{(\mathbf{E}\mathbf{e})}_{\mathbf{0}}\mathbf{e}^T = -\mathbf{E} \quad (4)$$

Substituting eq.(4) into eq.(3), we obtain

$$-\mathbf{E}\mathbf{u} = \lambda^3\mathbf{u} \quad (5)$$

Moreover, from eq.(1) we have  $\mathbf{E}\mathbf{u} = \lambda\mathbf{u}$ , and thus eq.(5) reduces to

$$\lambda(1 + \lambda^2)\mathbf{u} = \mathbf{0} \quad (6)$$

However,  $\|\mathbf{u}\| = 1$  and hence,  $\mathbf{u} \neq \mathbf{0}$ . Therefore, we have from eq.(6)

$$P(\lambda) \equiv \lambda(1 + \lambda^2) = 0 \quad (7)$$

Equation (7) is the characteristic equation of  $\mathbf{E}$ , and  $P(\lambda)$  its characteristic polynomial. The three roots of eq.(7) are thus

$$\lambda_1 = 0, \quad \lambda_2 = j, \quad \lambda_3 = -j$$

The eigenvector  $\mathbf{u}_1$  corresponding to  $\lambda_1$  is, obviously,  $\mathbf{e}$ , for  $\mathbf{E}\mathbf{e} = \mathbf{0}$ . Now, from eqs.(2) and (3), note that  $\mathbf{E}$  and its integer powers share the same set of eigenvectors.

Now, to find  $\mathbf{u}_2$ , corresponding to  $\lambda_2$ , let us substitute these into eq.(2), along with  $\mathbf{E}^2 = -\mathbf{1} + \mathbf{e}\mathbf{e}^T$ ,

$$(-\mathbf{1} + \mathbf{e}\mathbf{e}^T)\mathbf{u}_2 = -\mathbf{u}_2 \quad \text{or} \quad \mathbf{u}_2 + (\mathbf{e}^T\mathbf{u}_2)\mathbf{e} = -\mathbf{u}_2$$

which implies that

$$\mathbf{e}^T\mathbf{u}_2 = 0$$

Thus,  $\mathbf{u}_2$  is any real vector normal to  $\mathbf{e}$ , which means that  $\mathbf{u}_2$  lies in the plane normal to  $\mathbf{e}$ . Moreover, since  $\mathbf{E}^2$  is symmetric, its three eigenvectors are mutually orthogonal,  $\mathbf{u}_3$  thus being, simply,

$$\mathbf{u}_3 = \mathbf{e} \times \mathbf{u}_2$$

In conclusion,  $\mathbf{u}_2$  and  $\mathbf{u}_3$  are any two mutually orthogonal unit vectors lying in the complex plane perpendicular to  $\mathbf{e}$ , e.g.

$$\mathbf{u}_2 = [0 \quad \alpha \quad \beta]^T$$

where  $\alpha$  and  $\beta$  are two complex numbers, so that

$$\|\mathbf{u}_2\|^2 \equiv \mathbf{u}_2^*\mathbf{u}_2 = [0 \quad \bar{\alpha} \quad \bar{\beta}] \begin{bmatrix} 0 \\ \alpha \\ \beta \end{bmatrix} = \bar{\alpha}\alpha + \bar{\beta}\beta = |\alpha|^2 + |\beta|^2 = 1$$

where  $(\cdot)^*$  is the *transpose conjugate* of the complex vector  $(\cdot)$ ,  $(\bar{\cdot})$  denotes conjugate of  $(\cdot)$ , and  $|\cdot|$  the module of  $(\cdot)$ . Thus

$$\mathbf{u}_2^*\mathbf{u}_3 = 0$$