Questions

1. Consider the following sequence of stack operations:
   \text{push(d), push(h), pop(), push(f), push(s), pop(), pop(), push(m).}
   (a) Assume the stack is initially empty, what is the sequence of popped values, and what is the final state of the stack? (Identify which end is the top of the stack.)
   (b) Suppose you were to replace the \text{push} and \text{pop} operations with \text{enqueue} and \text{dequeue} respectively. What would be the sequence of dequeued values, and what would be the final state of the queue? (Identify which end is the front of the queue.)

2. Use a stack to test for balanced parentheses, when scanning the following expressions. Your solution should show the state of the stack each time it is modified. The “state of the stack” must indicate which is the top element.
   Only consider the parentheses \([,],(,),\{,\}. Ignore the variables and operators.
   (a) \([ a + \{ b / ( c - d ) + e / (f + g ) } - h ]
   (b) \([ a \{ b + [ c ( d + e ) - f ] + g \}

3. Suppose you have a stack in which the values 1 through 5 must be pushed on the stack in that order, but that an item on the stack can be popped at any time. Give a sequence of push and pop operations such that the values are popped in the following order:
   (a) 2, 4, 5, 3, 1
   (b) 1, 5, 4, 2, 3
   (c) 1, 3, 5, 4, 2
   It might not be possible in each case.

4. (a) Suppose you have three stacks \text{s1}, \text{s2}, \text{s2} with starting configuration shown on the left, and finishing condition shown on the right. Give a sequence of push and pop operations that take you from start to finish. For example, to pop the top element of \text{s1} and push it onto \text{s3}, you would write \text{s3.push( s1.pop()).}

<table>
<thead>
<tr>
<th>start</th>
<th>finish</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>A</td>
</tr>
<tr>
<td>B</td>
<td>B</td>
</tr>
<tr>
<td>C</td>
<td>D</td>
</tr>
<tr>
<td>D</td>
<td>C</td>
</tr>
<tr>
<td>s1</td>
<td>s1</td>
</tr>
<tr>
<td>s2</td>
<td>s2</td>
</tr>
<tr>
<td>s3</td>
<td>s3</td>
</tr>
</tbody>
</table>

last updated: 2\textsuperscript{nd} Oct, 2016 at 06:37
(b) Same question, but now suppose the finish configuration on s3 is BDAC (with B on top)?

5. Consider the following sequence of stack commands:

\text{push}(a), \text{push}(b), \text{push}(c), \text{pop}(), \text{push}(d), \text{push}(e), \text{pop}(), \text{pop}(), \text{pop}(), \text{pop}().

(a) What is the order in which the elements are popped? (Give a list and indicate which was popped first.)

(b) Change the position of the \text{pop}() commands in the above sequence so that the items are popped in the following order: \text{b,d,c,a,e}.

You are not allowed to change the ordering of the \text{push} commands.

I briefly discussed the next two questions in the lecture (slides only)

6. Assume you have a stack with operations: \text{push}(), \text{pop}(), \text{isEmpty}(). How would you use these stack operations to simulate a queue, in particular, the operations \text{enqueue}() and \text{dequeue}()?

Hints: Use two stacks one of which is the main stack and one is a temporary one. You are allowed use a \text{while} loop.

If you have no idea how to do this question or what is even being asked, then look at the solutions. Once you understand the solutions, try to do the next question and resist looking at the solution for that one.

7. Assume you have a queue with operations: \text{enqueue}(), \text{dequeue}(), \text{isEmpty}(). How would you use the queue methods to simulate a stack, in particular, \text{push}() and \text{pop}()?

Hint: use two queues, one of which is the main one and one is temporary. The solution requires that at least one of the \text{push}() or \text{pop}() operations is \(O(N)\) where \(N\) is the size of the queue.
Answers

1. (a) Sequence of popped values: h,s,f. State of stack (from top to bottom): m, d
   (b) Sequence of dequeued values: d,h,f. State of queue (from front to back): s,m.

2. (a) - means empty stack

   ```
   [
   [{}
   {{}{( TOP OF STACK IS ON THE RIGHT
   [{
   [{}[{
   [{}[{
   [{}[{
   [{}[{
   [ empty stack, so brackets match
   ```

   (b) -

   ```
   [
   [{}[{
   [{}[{
   [{}[{
   [{}[{
   [ stack not empty, so brackets don’t match
   ```

3.

<table>
<thead>
<tr>
<th>24531</th>
<th>15423</th>
<th>13542</th>
</tr>
</thead>
<tbody>
<tr>
<td>push 1</td>
<td>push 1</td>
<td>push 1</td>
</tr>
<tr>
<td>push 2</td>
<td>pop</td>
<td>pop</td>
</tr>
<tr>
<td>pop</td>
<td>push 2</td>
<td>push 2</td>
</tr>
<tr>
<td>push 3</td>
<td>push 3</td>
<td>push 3</td>
</tr>
<tr>
<td>push 4</td>
<td>push 4</td>
<td>pop</td>
</tr>
<tr>
<td>pop</td>
<td>push 5</td>
<td>push 4</td>
</tr>
<tr>
<td>push 5</td>
<td>pop</td>
<td>push 5</td>
</tr>
<tr>
<td>pop</td>
<td>pop</td>
<td>pop</td>
</tr>
<tr>
<td>pop</td>
<td>x</td>
<td>pop</td>
</tr>
<tr>
<td>pop</td>
<td>(not possible)</td>
<td>pop</td>
</tr>
</tbody>
</table>
4. (a) s2.push( s1.pop() )
    s2.push( s1.pop() )
    s3.push( s1.pop() )
    s3.push( s1.pop() )
    s3.push( s2.pop() )
    s3.push( s2.pop() )

(b) s2.push( s1.pop() )
    s2.push( s1.pop() )
    s3.push( s1.pop() )
    s1.push( s2.pop() )
    s3.push( s2.pop() )
    s2.push( s1.pop() )
    s3.push( s1.pop() )
    s3.push( s2.pop() )

5. (a) c (popped first), e, d, b, a
    (b) push(a), push(b), pop(), push(c), push(d), pop(), pop(), push() push(e), pop()

6. Here I present two solutions. Note that each requires that at least one of the `enqueue()` or `dequeue()` operation is $O(N)$ where $N$ is the size of the queue.

**Solution 1**

The first solution is to implement `enqueue(e)` simply by pushing the new element onto the main stack $s$.

```java
enqueue(e){
    s.push(e)
}
```

In this solution, the bottom of the stack would be the front of the queue (containing the element that has been in the queue longest) and the top of the stack is the back of the queue (containing the least recently added element). How can we implement `dequeue`, that is, how do we remove the bottom element of the stack?

The idea is to use a second stack $tmpS$. We first pop all items from the main stack $s$, and push each popped element directly onto the second stack $tmpS$. We then pop the top element of the second stack (which is the oldest element in the set). Finally, we refill the main stack $s$ by popping all elements from the second stack and pushing each back on to the first stack.

```java
dehqueue(){
    tmpS <- new empty stack
    while !s.isEmpty(){
        tmpS.push( s.pop() )
    }
    s.pop()
    while !tmpS.isEmpty(){
        s.push( tmpS.pop() )
    }
}
```
returnValue <- tmpS.pop() // the dequeued element
while !(tmpS.isEmpty()){
    s.push( tmpS.pop() )
}
return returnValue
}

Solution 2

Here we let dequeue() be simple and just pop the stack. For this to work, the stack needs to store the elements such that the oldest element is on top of the stack and the most recently added element is at the bottom of the stack.

dequeue(){
    return s.pop()
}

With this solution, enqueue(e) needs to do the heavy lifting: enqueue uses a temporary stack to invert the order of elements currently in the main stack, so that the newest element is on top of this temporary stack. Then it pushes the new element onto the empty main stack. Then it copies all the elements back from the temporary stack to the main stack, using pop-push. This recreates the original stack, but now the newest element has been added on the bottom.

enqueue(e){
    tmp <- new empty stack
    while ! (s.isEmpty()){
        tmpS.push( s.pop() )
    }
    s.push(e)
    while ! (tmpS.isEmpty()){
        s.push( tmpS.pop() )
    }
}
7. The concepts are similar to the previous question.

**Solution 1**

```java
push(e){
    q.enqueue(e)
}

drop(){
    tmpQ <- new empty queue
    while ! (q.isEmpty()){  
        tmpE <- q.dequeue()
        if ! (q.isEmpty())
            tmpQ.enqueue( tmpE )
        else{
            while !( tmpQ.isEmpty())
                q.enqueue( tmpQ.dequeue())
            return tmpE
        }
    }
}
```

**Solution 2**

Here the idea is similar, but now push does most of the work.

```java
drop(){
    q.dequeue()
}

push(e){
    make a new empty queue qTmp
    qTmp.enqueue(e)
    while !q.isEmpty(){
        qTmp.enqueue( q.dequeue() )
    }
    q <- qTmp
    return qTmp
}
```