

COMP 250

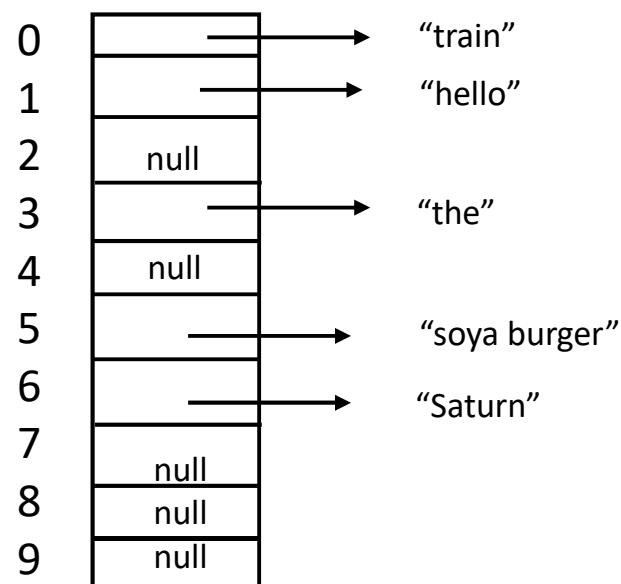
Lecture 9

Array lists

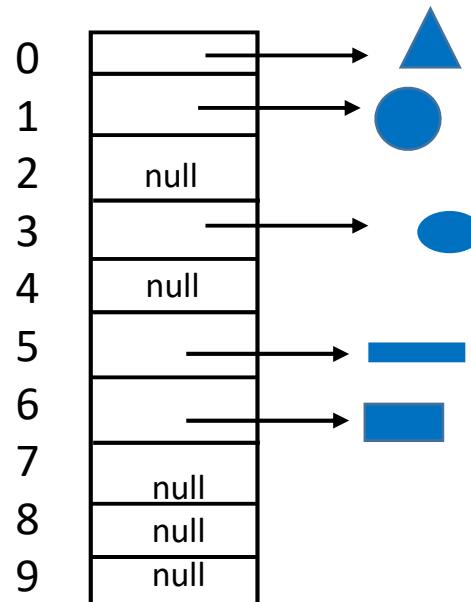
Jan. 26, 2022

Arrays of reference type variables

String []



Shape []



```
Shape[] shapes = new Shape[428];
shapes[3] = new Shape(      );
// etc...
```

Arrays versus ‘Array Lists’

Arrays contain elements. There is no restriction on *where* the elements are stored.

Elements can be accessed in “constant time”.

Arrays *can be used* to make lists, sometimes called ‘array lists’.

Java has an `ArrayList` class.

List

A list is an ordered set of elements

$$a_0, a_1, a_2, a_3, \dots, a_{N-1}$$

N is the number of elements in the list, usually called the “size” of the list.

What things do we do with a list?

- get(i)** // Returns the i-th element (but doesn't remove it)
- set(i,e)** // Replaces the i-th element with e
- add(i,e)** // Inserts element e into the i-th position
- remove(i)** // Removes the i-th element from list
- remove(e)** // Removes first (if any) occurrence of element e
- clear()** // Empties the list.
- isEmpty()** // Returns true if empty, false if not empty.
- size()** // Returns number of elements in the list

.....

Lists

- **arraylist (today)**

- singly linked list

- doubly linked list

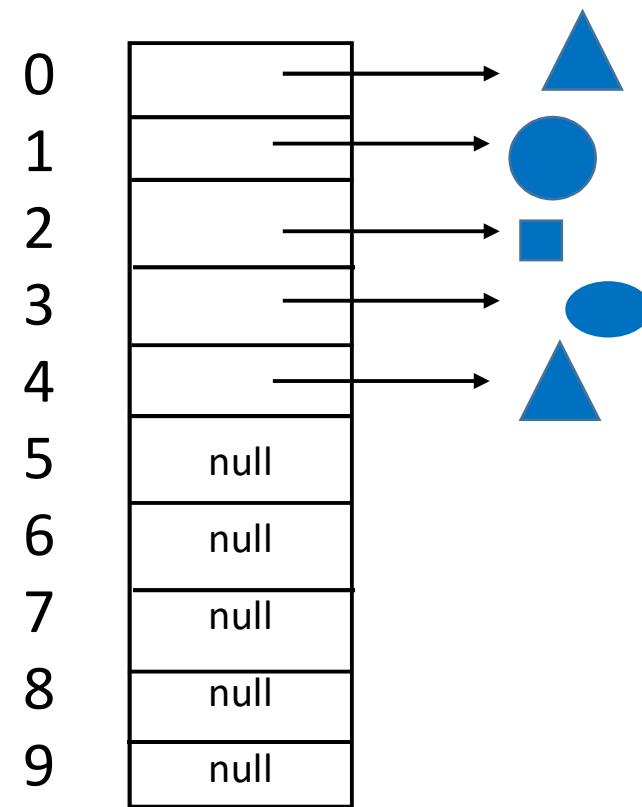
- :



next two lectures

size = 5
length = 10

array list of Shape



How to implement various operations ?

(pseudocode only -- no return type given)

get(i)

set(i,e)

add(i,e)

remove(i)

remove(e)

clear()

isEmpty()

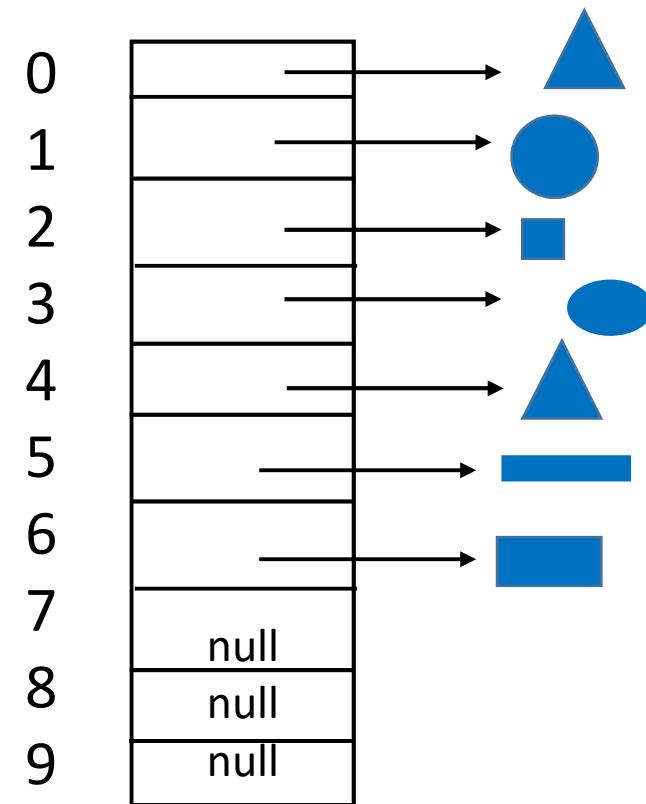
size()

.....

```
get(i) {
```

```
}
```

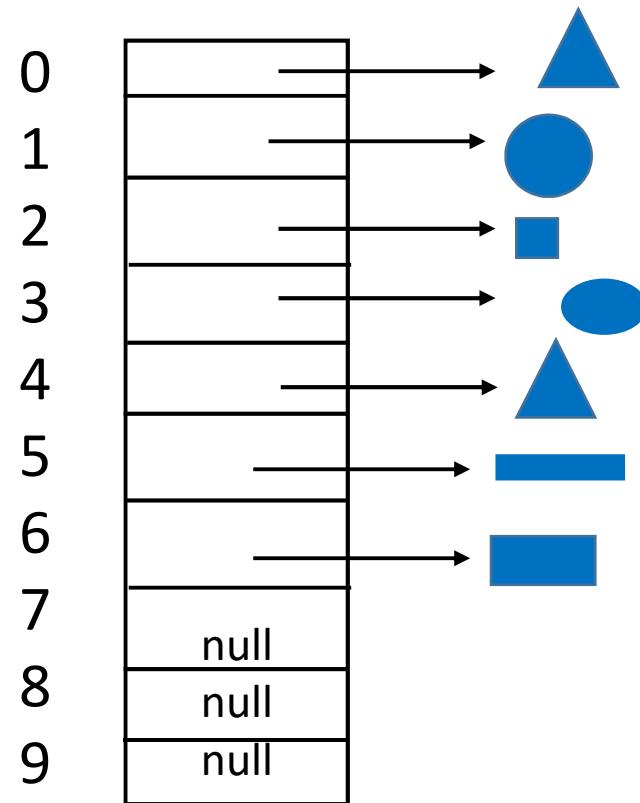
```
// Let a[ ] be the array.
```



size = 7

length = 10

```
get(i) {  
    if (i >= 0) & (i < size)  
        return a[ i ]  
}
```



size = 7
length = 10

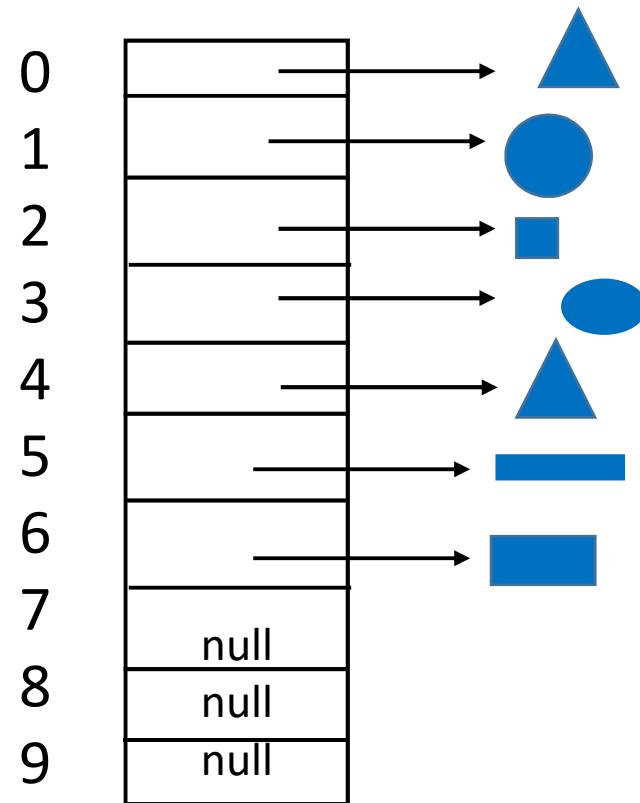
```

get(i) {

    if (i >= 0) & (i < size)
        return a[ i ]
    else
        // throw an index
        // out of bounds
        // exception
}

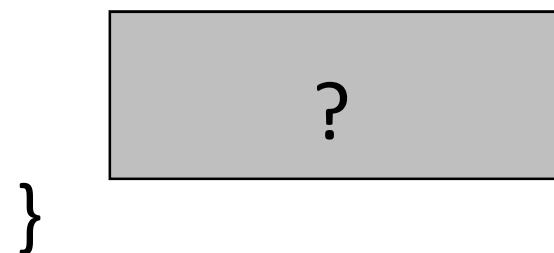
```

I will not mention this check in the rest of the methods today. But be aware that it would be added in proper implementation.

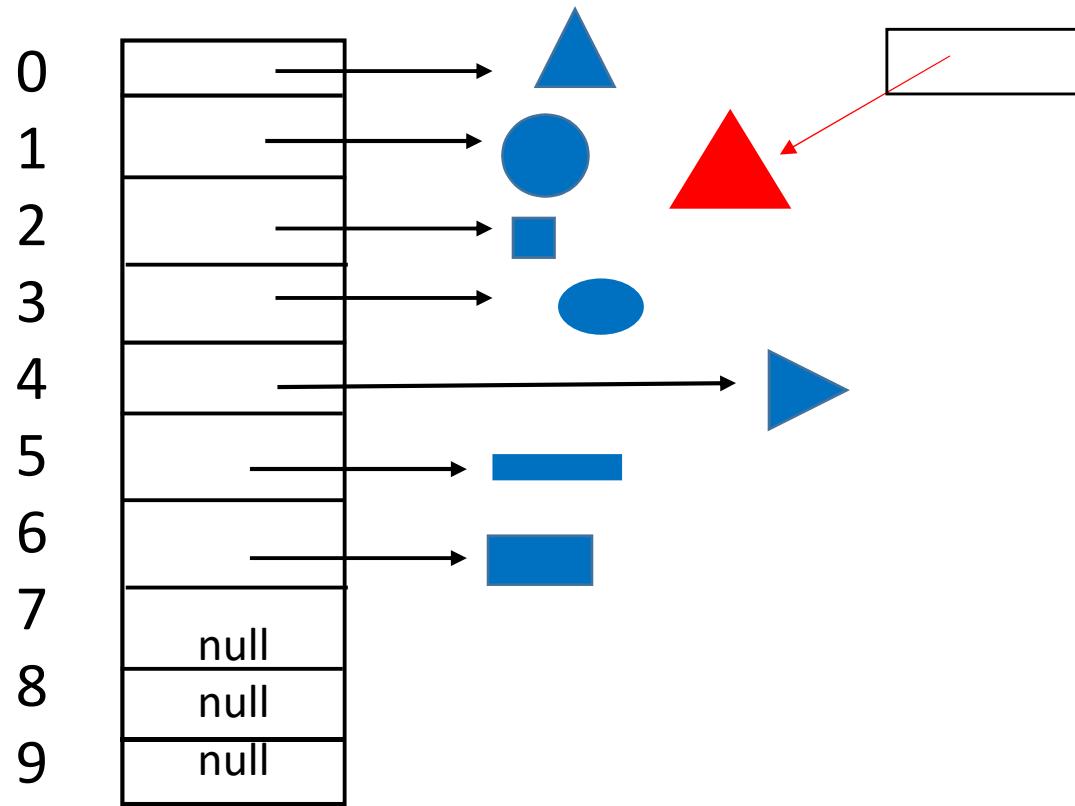


size = 7
length = 10

```
set(i,e){ // replaces the object at index i with object e
```



e.g. set(4, e)



I won't draw this box in the rest of the slides

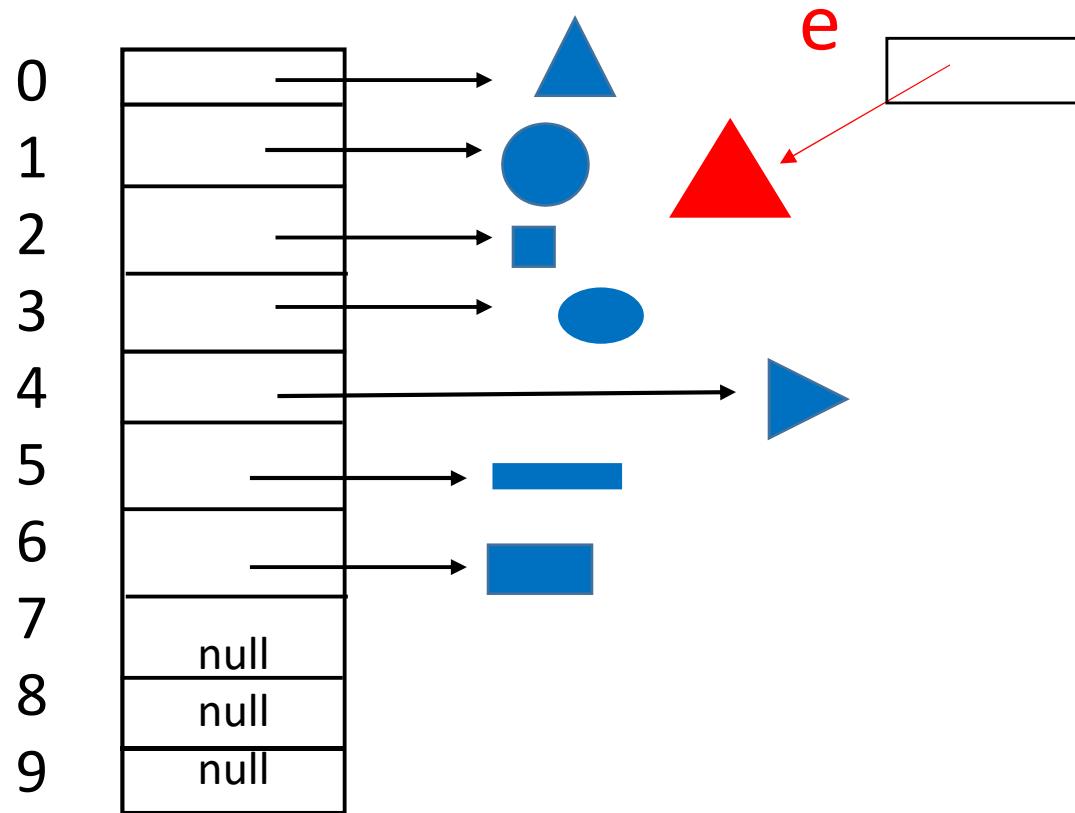
```
set(i,e){ // replaces the object at index i with object e
```

```
    if (i >= 0) & (i < size)
```

```
        a[i] = e
```

```
}
```

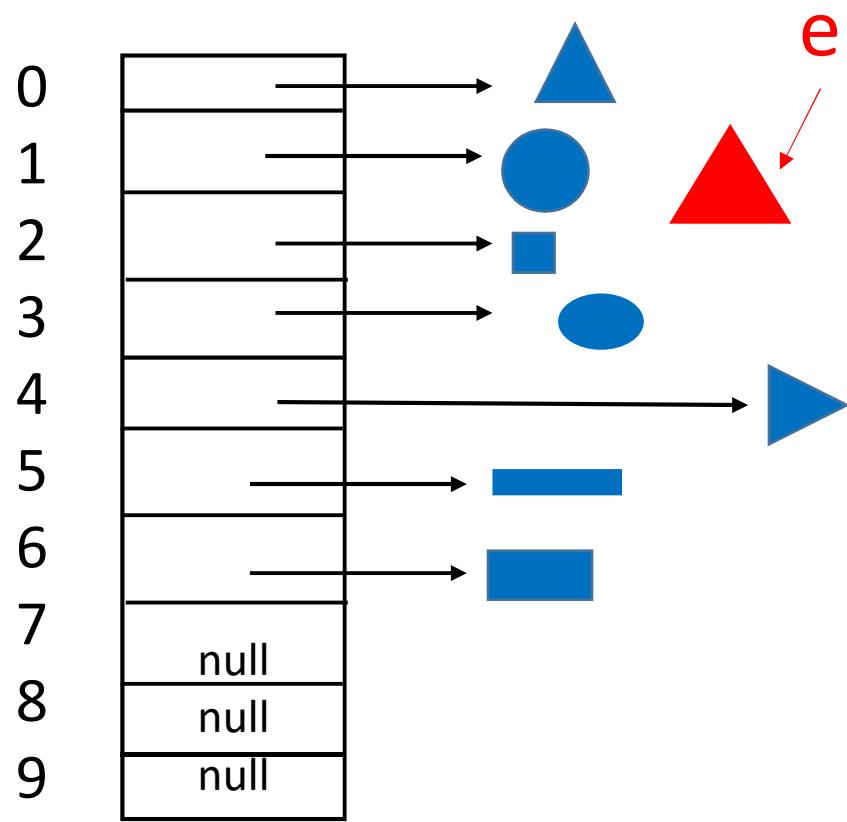
e.g. set(4, e)



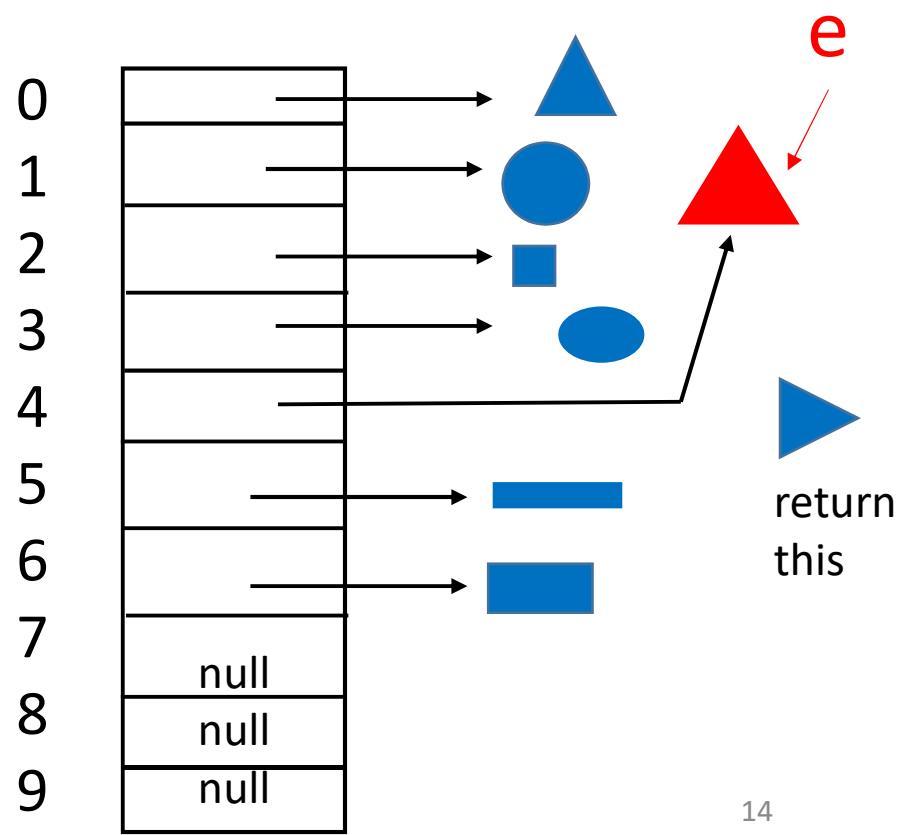
I won't draw this box in the rest of the slides

`set(4 ,e)`

BEFORE



AFTER



Alternatively, ...

```
set(i,e){    // replace the object at index i
             // and return the element that is replaced.

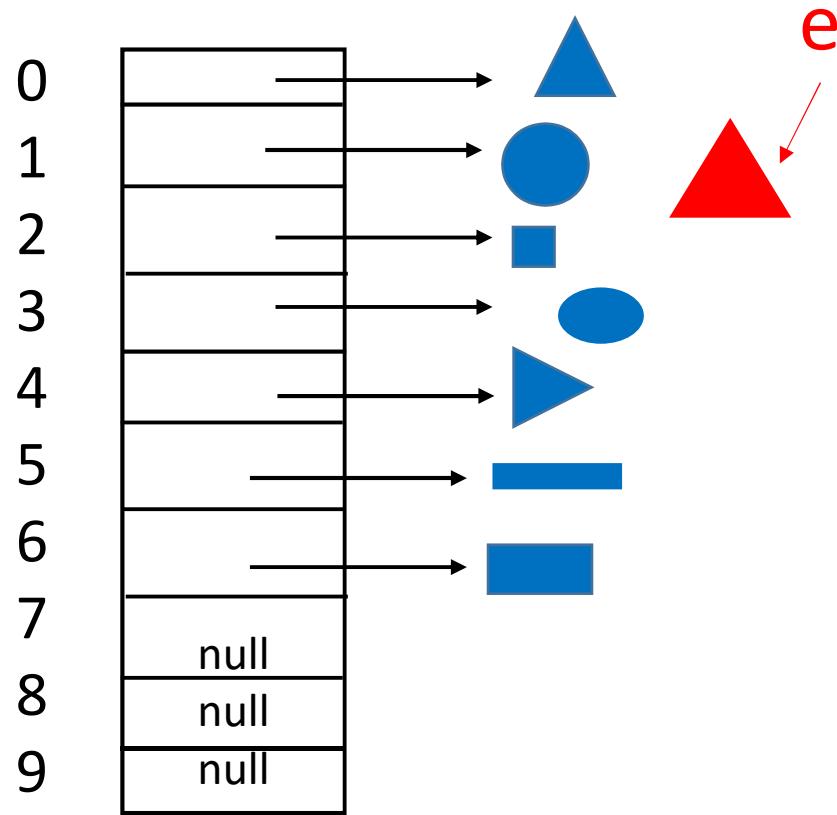
    if (i >= 0) & (i < size) {
        tmp = a[i]
        a[i] = e
    }
    return tmp
}
```

`add(i, e)`

// also known as “insert”

Make room by shifting the references, and then add reference to `e`

e.g. `add(2, e)`

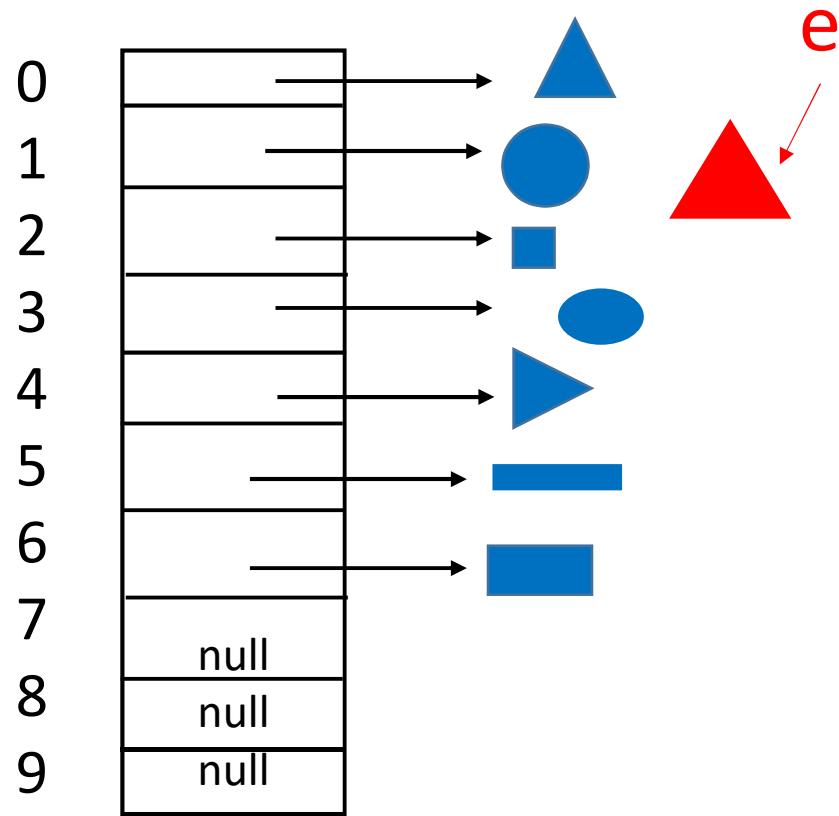


`add(2, e)`

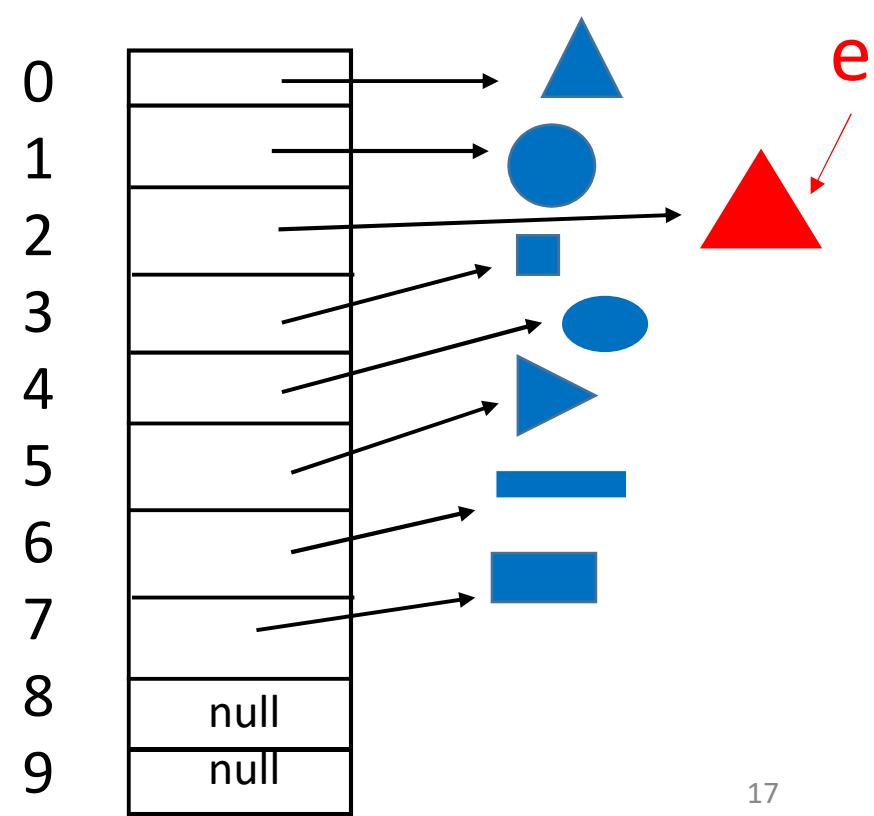
// also known as “insert”

Make room by shifting the references, and then add reference to `e`

BEFORE



AFTER



`add(i, e)`

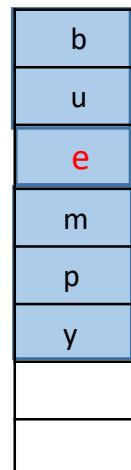
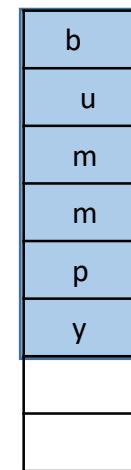
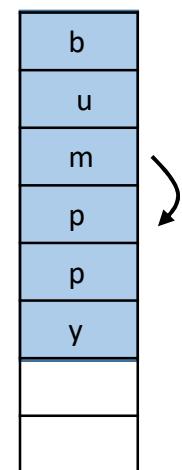
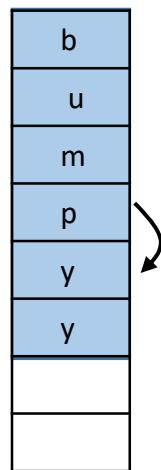
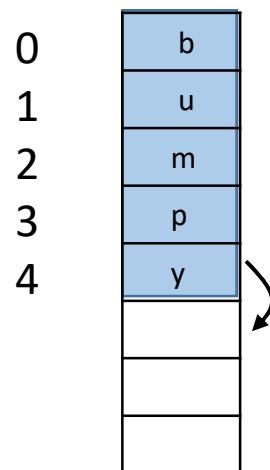
Heads up. I am using `e` both for parameter name and value here.

e.g. `add(2, e)`

Similar example, but now I store characters instead of shapes.

Insert character `e` into slot 2 of array list of characters: `bumpy`.

How? Make room by shifting. Then copy in the letter `e`.
This yields the list: `buempy`.



```
add( i, e) {
```

// in the figure below, i = 2, size = 5

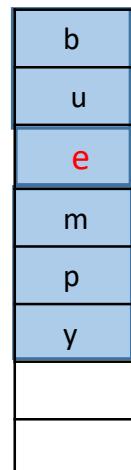
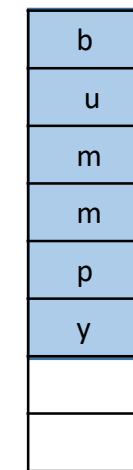
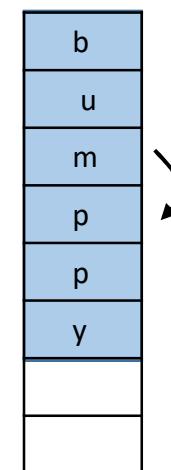
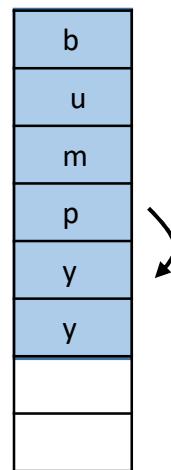
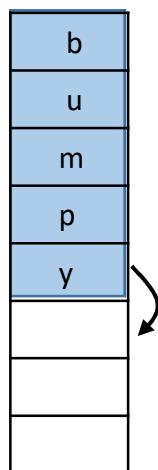
```
if (i >=0) & (i <= size) {
```

// todo: deal with case that size == length

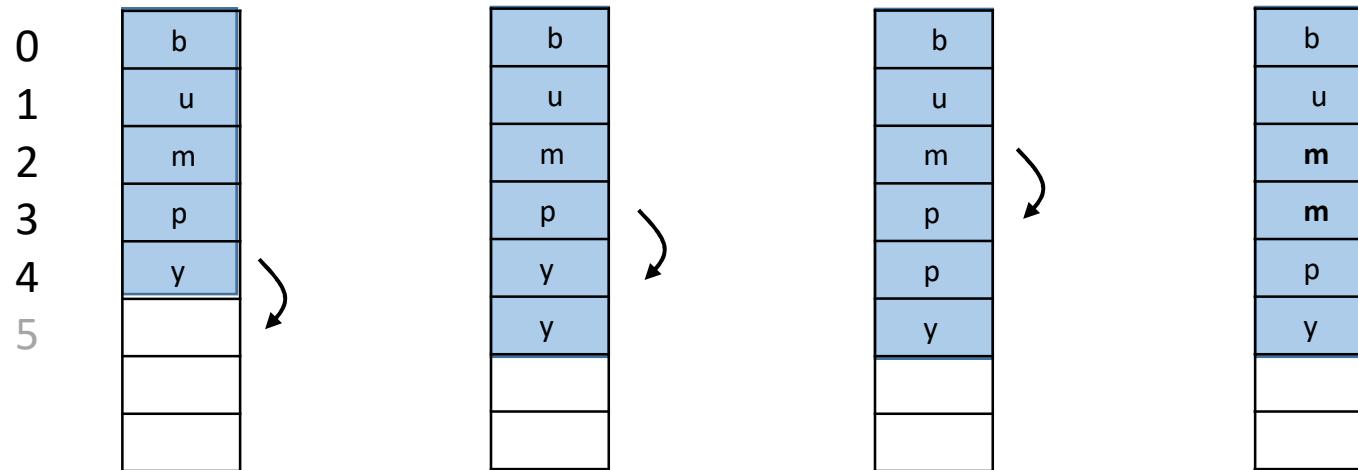
recall forward shifting pseudocode from lecture 4

```
}
```

```
}
```



```
add( i, e) {                                // in the figure below, i = 2, size = 5  
    if (i >= 0) & (i <= size) {  
        for (k = size; k > i; k--)  
            a[k] = a[k-1]           // sequence of copy forwards  
    }  
}
```



```
add( i, e) {
```

// in the figure below, i = 2

```
if (i >= 0) & (i <= size) {  
    for (k = size; k > i; k--)  
        a[k] = a[k-1]
```

a[i] = e

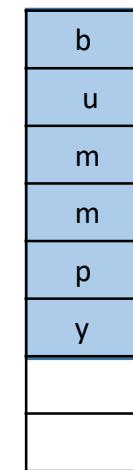
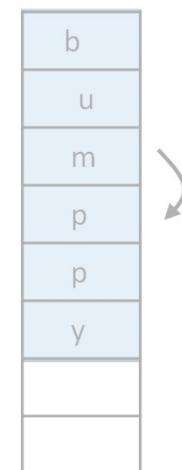
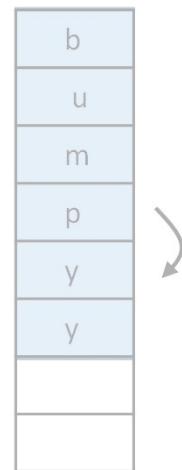
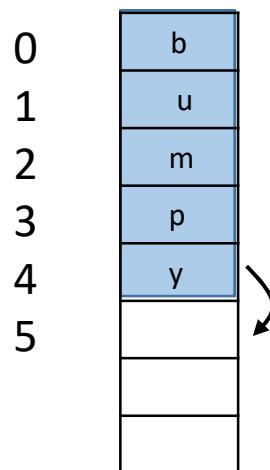
size = size + 1

}

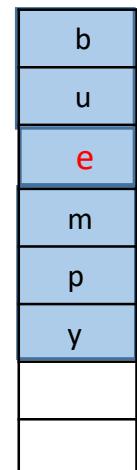
}

BEFORE

size = 5



AFTER
size = 6



How to add an element to an array list when array is full
(size == length) ?

```
add( i, e ) {  
  
    if (a.size == a.length){          // is array full?  
        make new bigger array b      // e.g. b.length = 1.5*a.length  
        for ( i=0; i < size; i++)  
            b[i] = a[i]                // copy elements to b  
  
        a = b  
    }  
  
    // insert here the add( i, e ) code from last slide  
}
```

Suppose you want to add an element to the list and you don't care where it goes?

Use `add(e)`, which calls `add(size, e)`. This adds to the end of the list.

If $\text{size} < \text{length}$, then this requires no shifting. It is therefore done in “constant time” (independent of number of elements in the list).

List Operations (ArrayList)

get(i)

set(i,e)

add(i,e)

remove(i) // Removes the i-th element from list

remove(e) // Removes element e from the list (if it is there)

clear() // Empties the list.

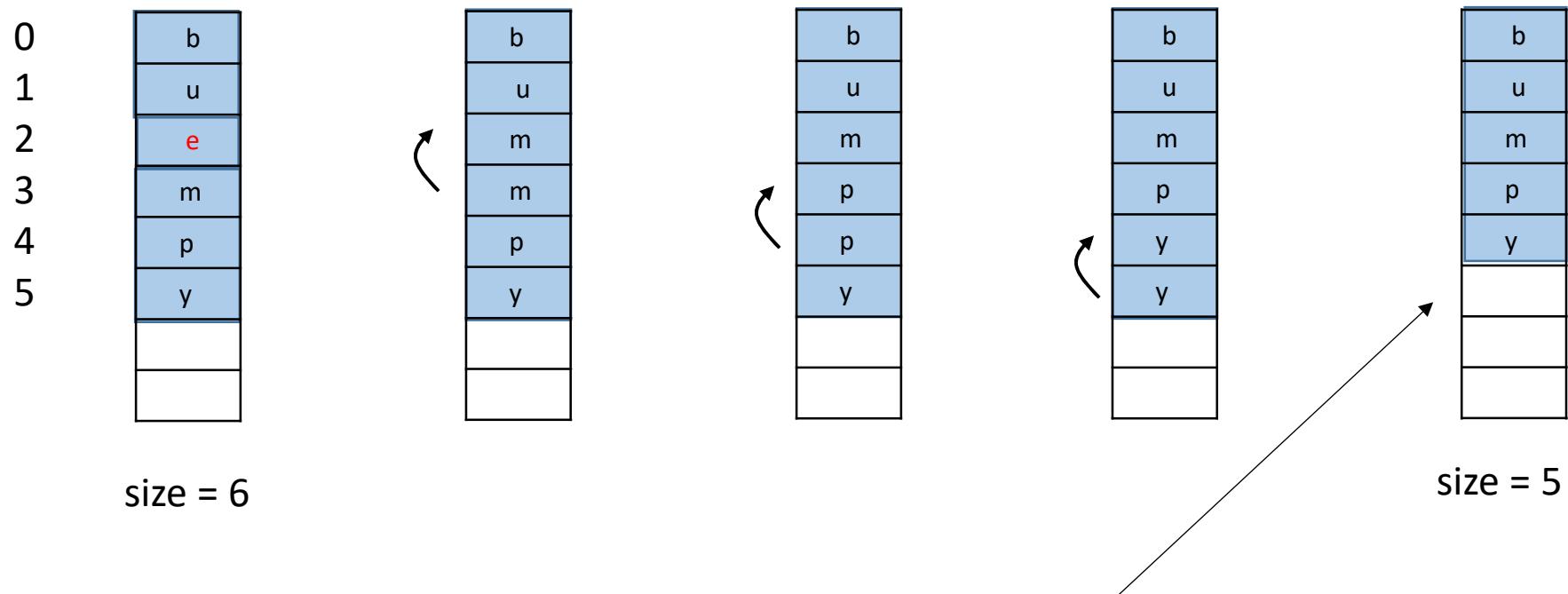
isEmpty() // Returns true if empty, false if not empty.

size() // Returns number of elements in the list

:

```
remove( i )           // e.g. i = 2 (removes the letter 'e')
```

It does the opposite of the add(i, **e**) method.



The value 'y' will still be there, but it is not accessible. Why not?

remove(i)

```
if ( (i >= 0) and (i < size) ){
```

```
    tmp = a[i]           // put aside and later return it
```

```
        Shift backwards the elements from position i+1 to end
```

```
    size = size - 1
```

```
    return tmp
```

```
}
```

remove(i)

```
if ( (i >= 0) and (i < size) ){

    tmp = a[i]                      // put aside and later return it

    for ( k = i; k < size - 1; k++){
        a[ k ] = a[ k + 1 ]
    }

    size = size - 1

    return tmp
}
```

Method Overloading

add(i ,e) // inserts element e into the i-th position

add(e) // inserts element e at end of list

remove(i) // Removes the i-th element from list

remove(e) // Removes first occurrence of element e
// from the list (if it is there)

Java ArrayList class

<https://docs.oracle.com/javase/8/docs/api/java/util/ArrayList.html>

It uses an array as the underlying data structure.

When the array is full and a new element is added, it “grows” the array by 50% namely, it creates a new array that is bigger and copies the elements to this new array.

Java generic types

An array of what? `ArrayList< T >`

Example:

```
ArrayList< Shape >   shape = new ArrayList< Shape >();
```

// initializes the array length (capacity) to 10

```
ArrayList< Shape >   shape = new ArrayList< Shape >( 23 );
```

// initializes the underlying array length to 23

Java generic types

An array of what? `ArrayList< T >`



It must be a reference type.

Example:

```
ArrayList< Integer > shape = new ArrayList< Integer >();
```

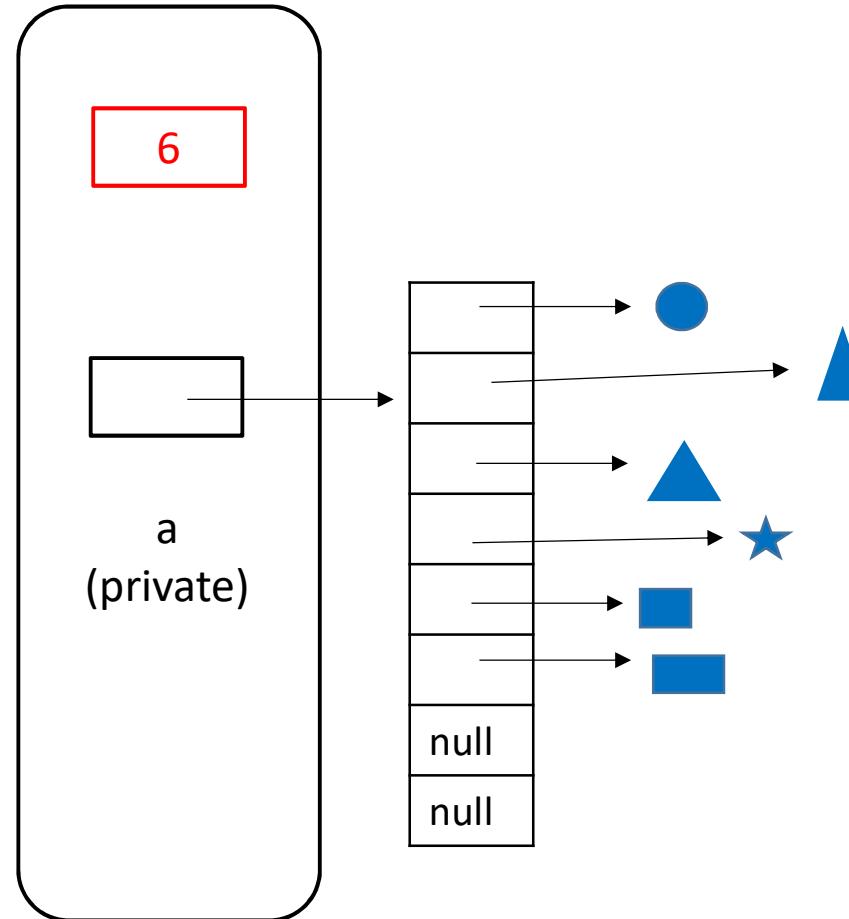
```
ArrayList< Integer > shape = new ArrayList< Integer >( 23 );
```

You cannot use a primitive type (`int`) in there.

Java ArrayList object

Has private field that holds the number of elements in the list (size).

Has a private field that references an array object.



How many objects are there in this figure? 8

Exercise (challenging)

Suppose you want to add N elements to an array list, which is initially empty.

Each time you fill the array, double its length.
(Java's ArrayList multiplies length by 1.5, not 2).

What is the $O()$ time complexity ?

Specifically...

Q: How many times k do we need to double the length of the array so that it can hold N elements ?

Q: How many **copy operations** are required ?
(When we increase the array length, we copy from the small array to large.)

See Exercises PDF for today. Try to answer the questions without peeking at the solutions.

Coming up...

Lectures

Fri. Jan. 28

Singly Linked Lists

Assessments

Fri. Jan. 28

Quiz 1 (lectures 1-7)

- practice quiz posted

Assignment 1 to be posted

- you will have 2 weeks to do it